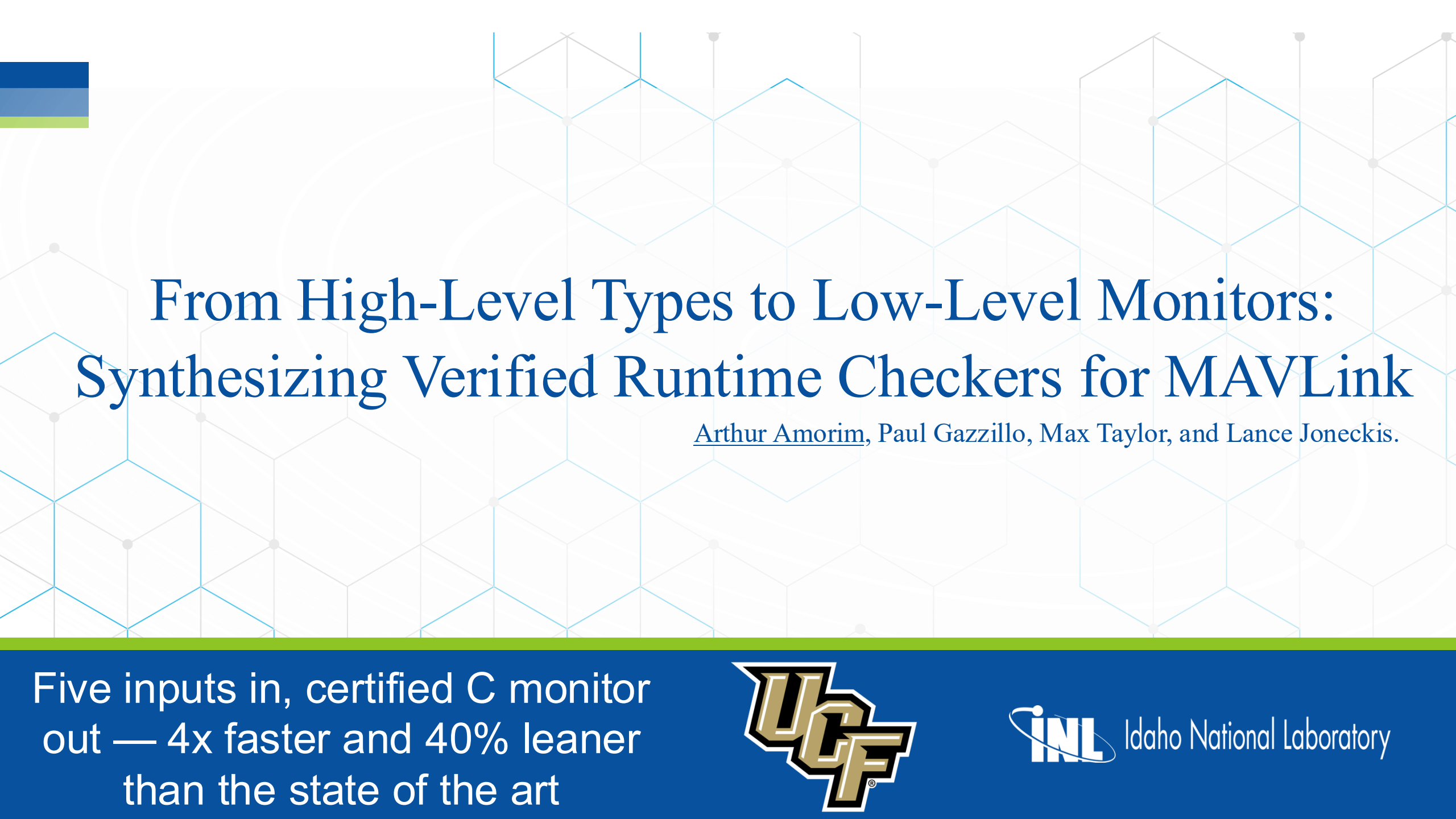




From High-Level Types to Low-Level Monitors: Synthesizing Verified Runtime Checkers for MAVLink

Arthur Amorim, Paul Gazzillo, Max Taylor, and Lance Joneckis.

Five inputs in, certified C monitor
out — 4x faster and 40% leaner
than the state of the art

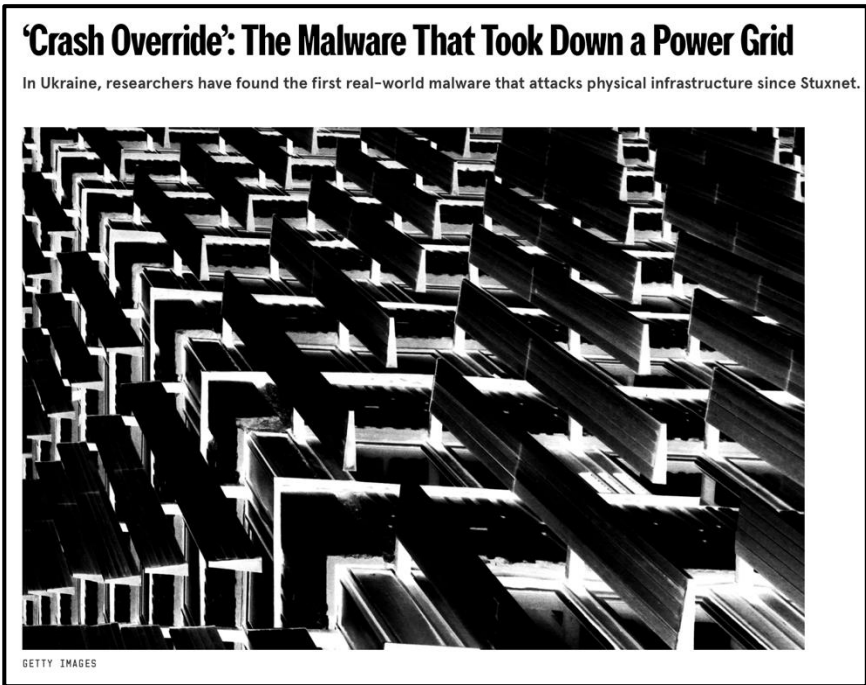


Cyber-Physical Systems Are Now Targets for Physical Damage via Cyber Attacks

Stuxnet



Crash Override



German Steel Mill



Stealthy Attacks Use Only Valid Commands in the Wrong Order

TECH + ENGINEERING

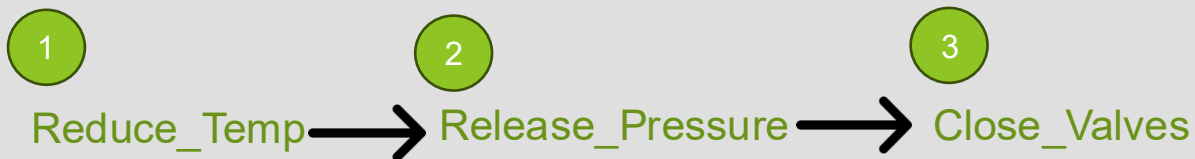
Cyber Attack on German Steel Mill Leads to 'Massive' Real World Damage

A steel mill in Germany lost control of its blast furnace. Hackers had infiltrated the mill's control system, according to the German government's office for information security.

BY R.A. BECKER THURSDAY, JANUARY 8, 2015 NOVA NEXT

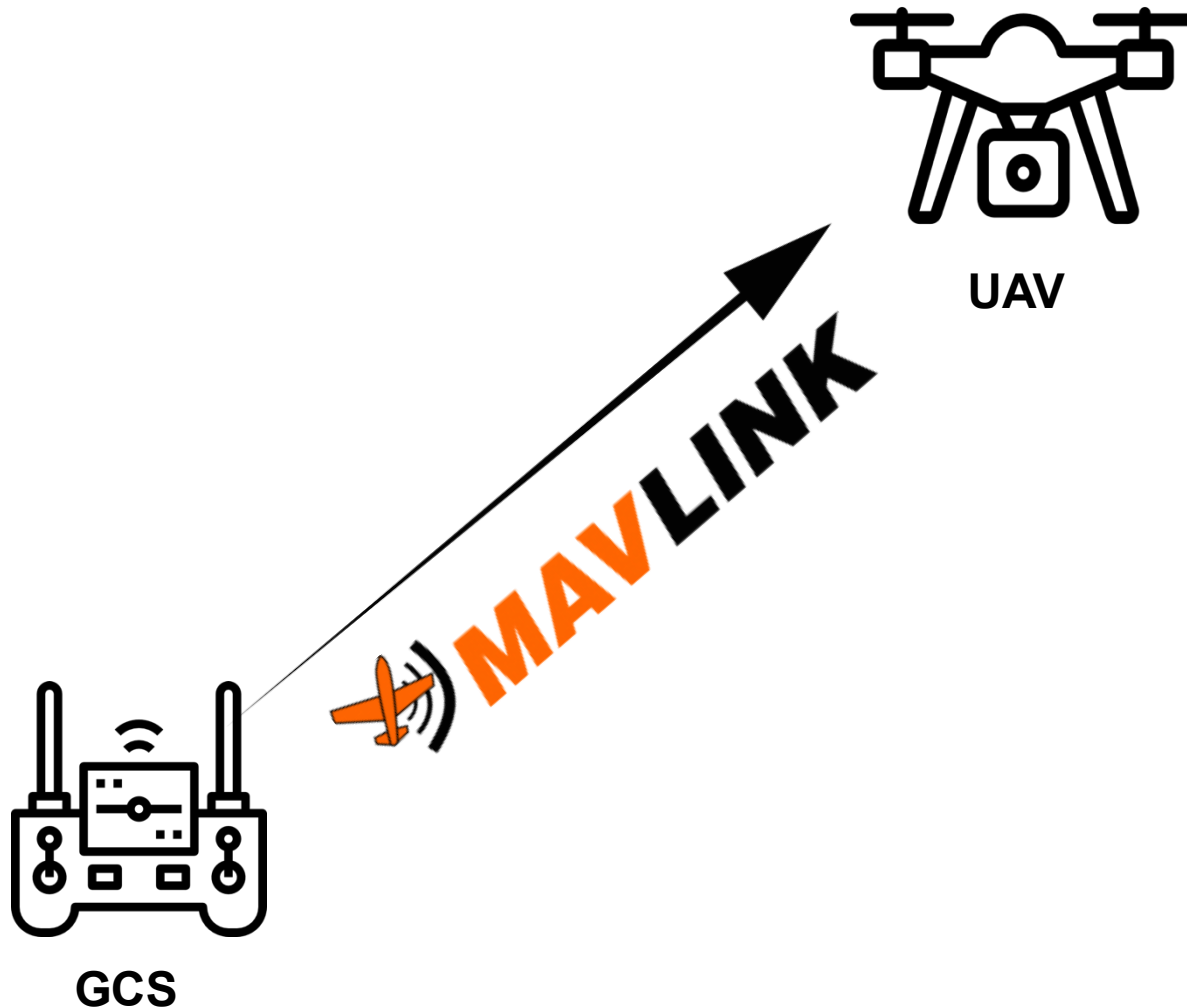


Blasting Furnace Shutdown Sequence:



Stealthy attack: valid commands, wrong order

Widely Used Protocols Are Susceptible to Stealthy Attacks



RVFUZZER: Finding Input Validation Bugs in Robotic Vehicles Through Control-Guided Testing

Taegyu Kim[†], Chung Hwan Kim^{*}, Junghwan Rhee^{*}, Fan Fei[†], Zhan Tu[†], Gregory Walkup[†]
Xiangyu Zhang[†], Xinyan Deng[†], Dongyan Xu[†]
[†]Purdue University, {tgkim, feif, tu17, gwalkup, xyzhang, xdeng, dxu}@purdue.edu
^{*}NEC Laboratories America, {chungkim, rhee}@nec-labs.com

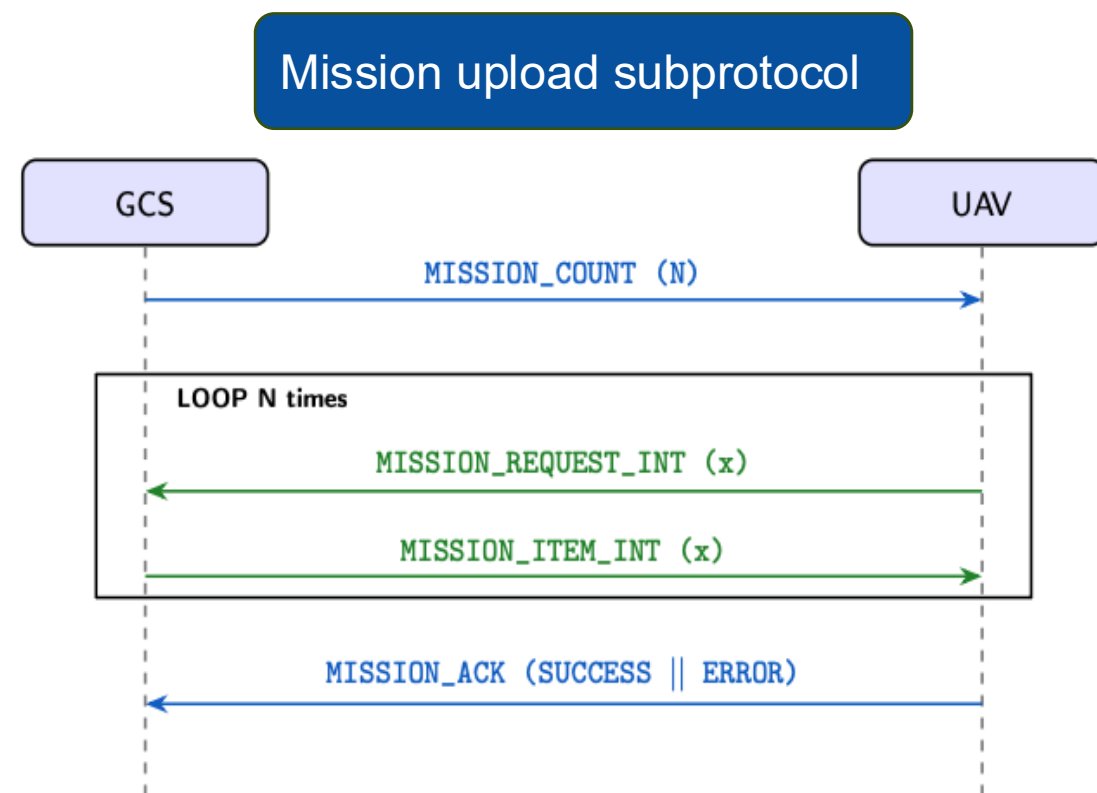
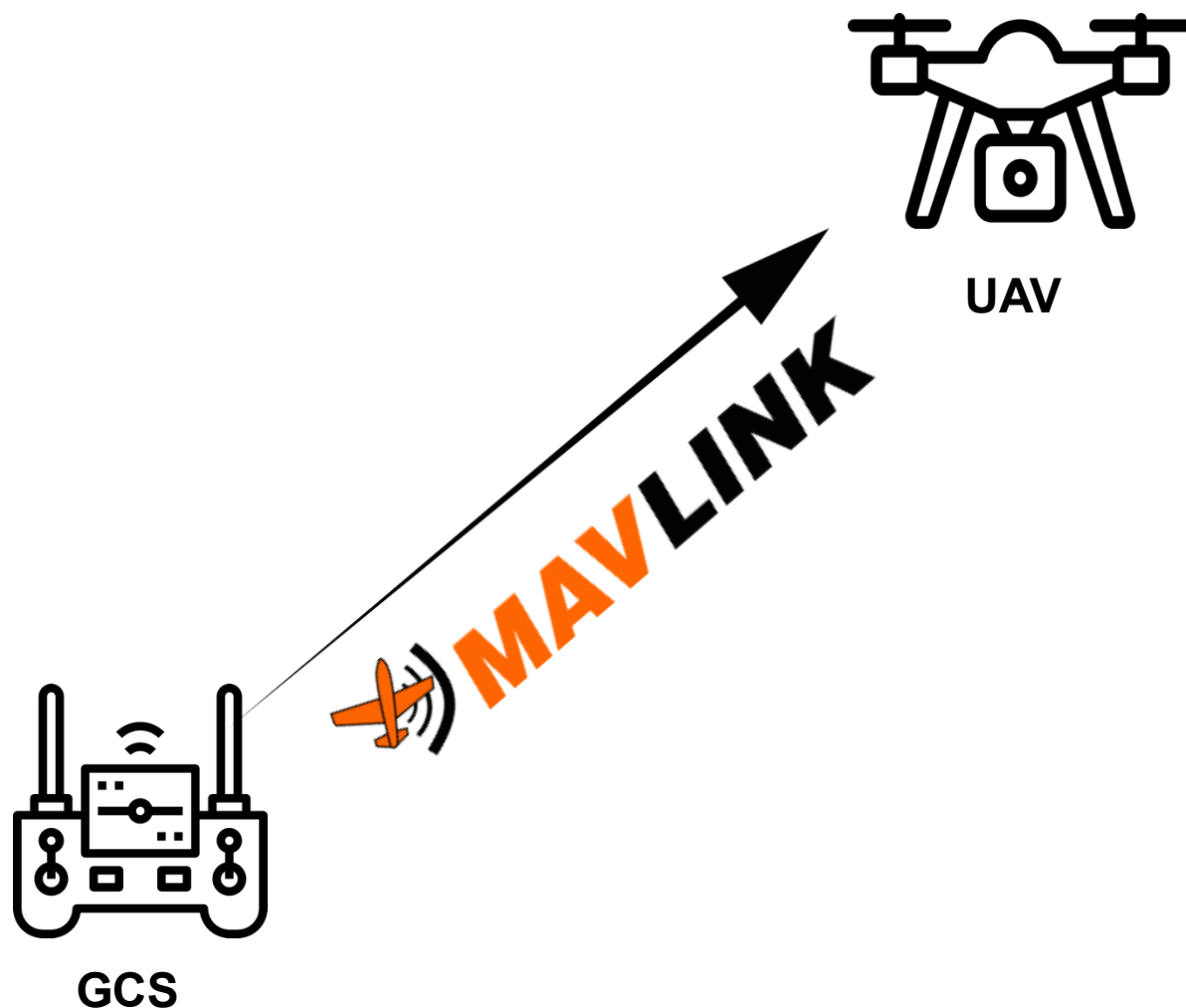
Found 89 input validation bugs

PGFUZZ: Policy-Guided Fuzzing for Robotic Vehicles

Hyungsub Kim, Muslum Ozgur Ozmen, Antonio Bianchi, Z. Berkay Celik, and Dongyan Xu
Purdue University
{kim2956, mozmen, antoniob, zcelik, dxu}@purdue.edu

Found 156 policy violation bugs

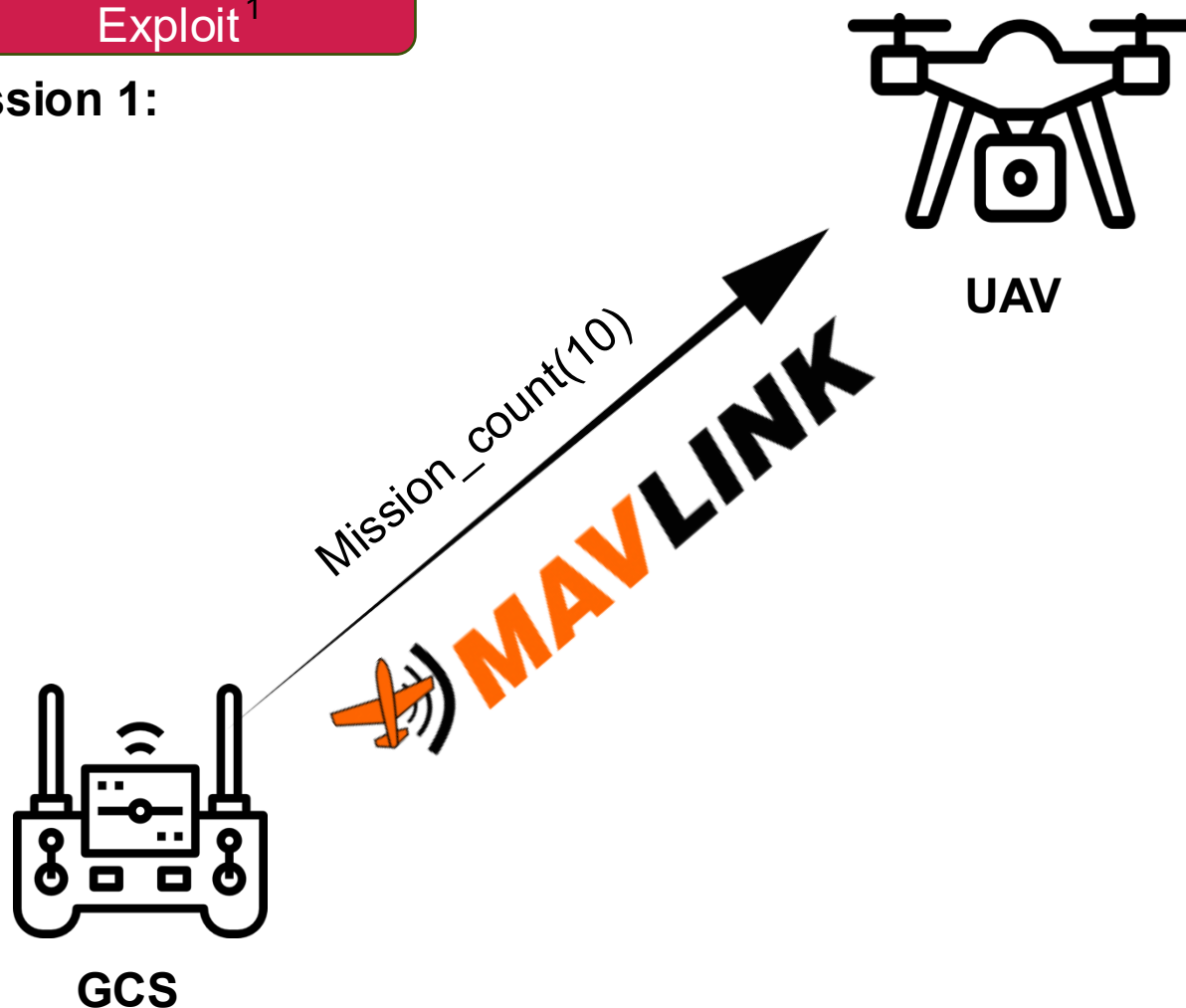
Widely Used Protocols Are Susceptible to Stealthy Attacks



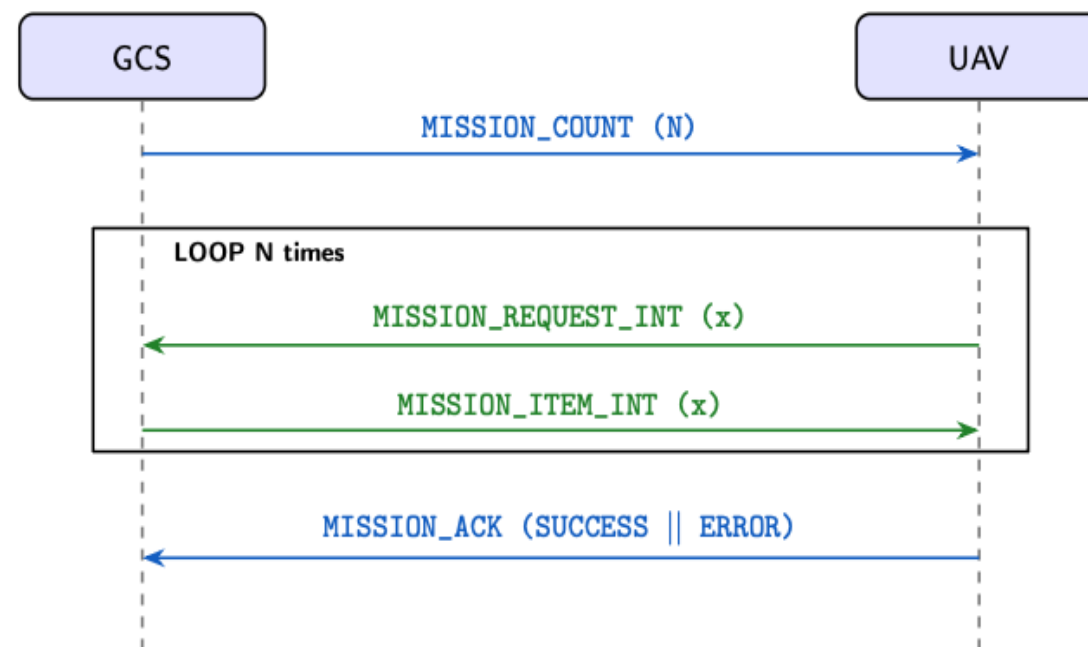
Widely Used Protocols Are Susceptible to Stealthy Attacks

Resource Usage
Exploit¹

Mission 1:



Mission upload subprotocol

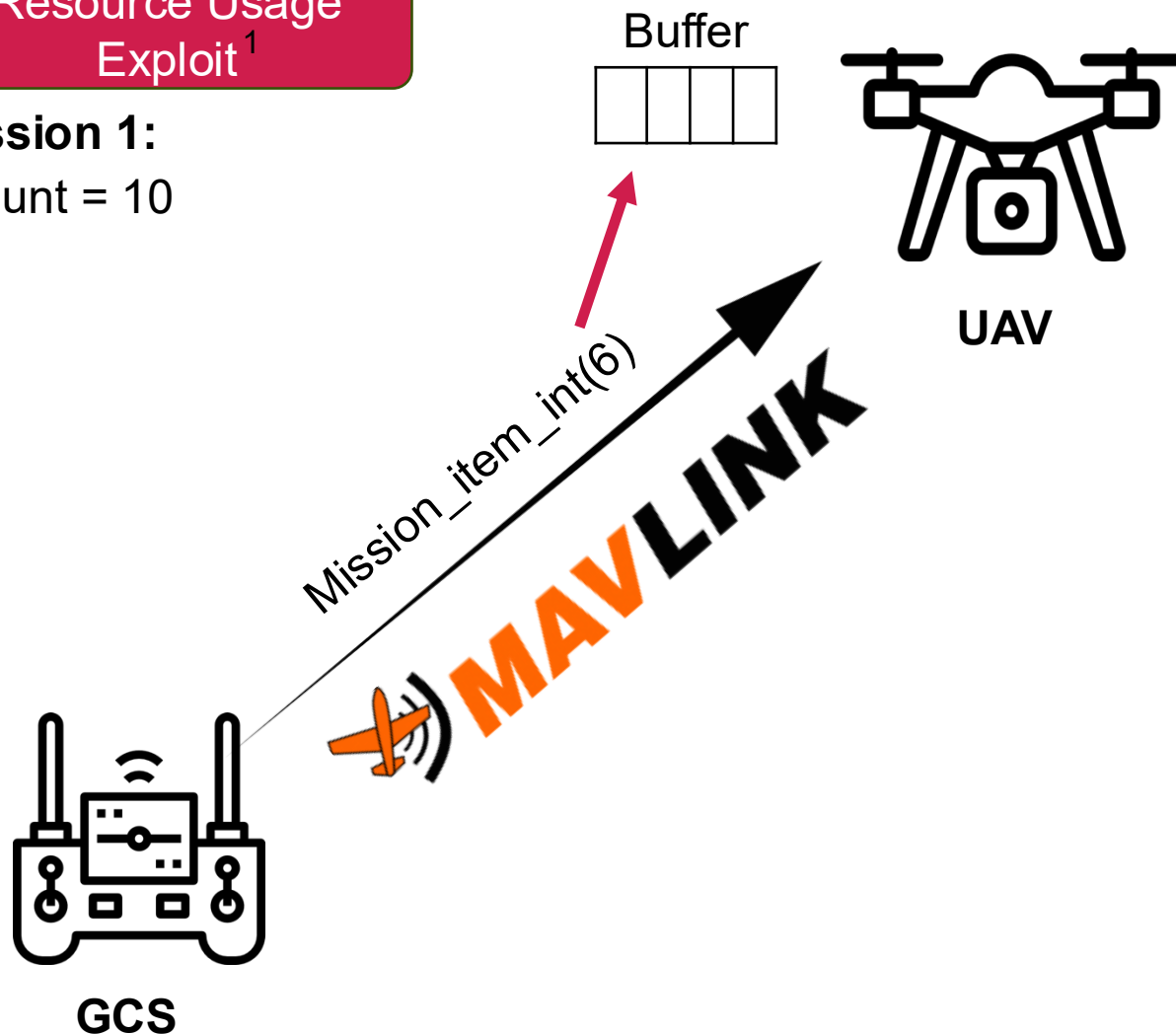


Widely Used Protocols Are Susceptible to Stealthy Attacks

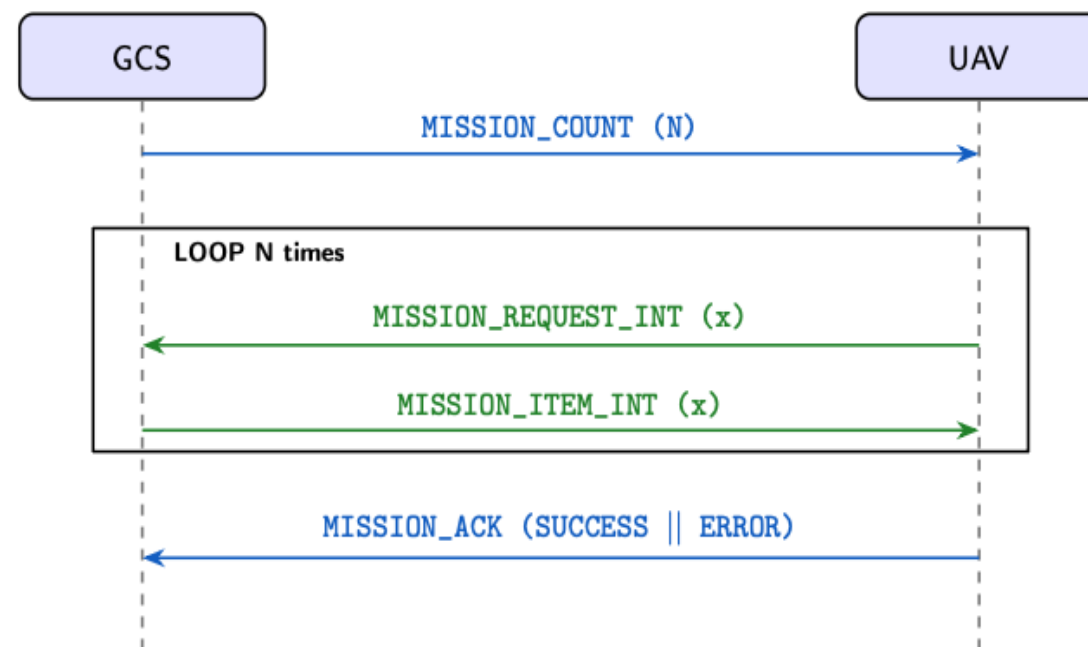
Resource Usage
Exploit¹

Mission 1:

Count = 10



Mission upload subprotocol



Widely Used Protocols Are Susceptible to Stealthy Attacks

Resource Usage
Exploit¹

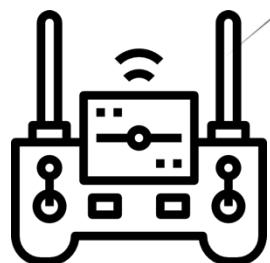
Mission 1:

Count = 10

Buffer
6 7 8 9



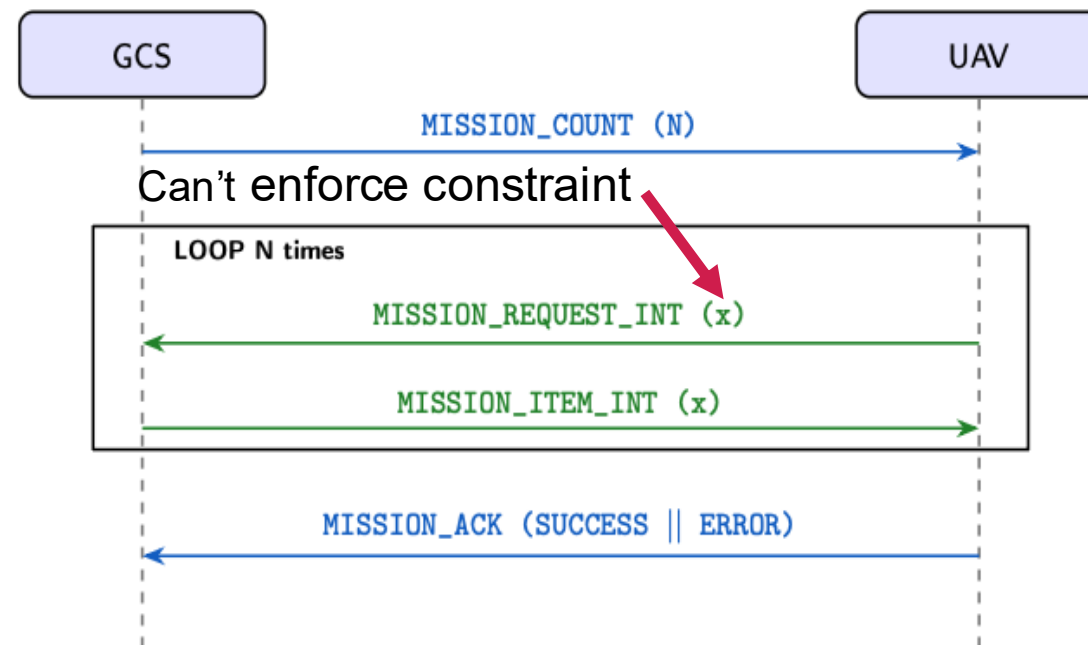
UAV



GCS

Mission_item_int(13)
MAV LINK

Mission upload subprotocol



Widely Used Protocols Are Susceptible to Stealthy Attacks

Resource Usage
Exploit¹

Mission 1:

Count = 10

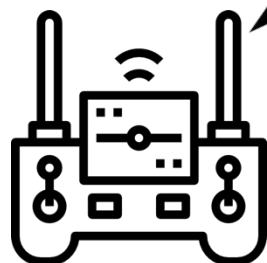
Buffer
6 7 8 9



UAV

Mission_ack(Success)

MAV LINK



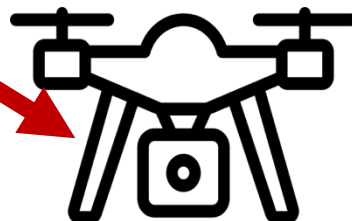
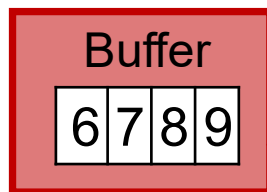
GCS

Widely Used Protocols Are Susceptible to Stealthy Attacks

Resource Usage
Exploit¹

Mission 2:

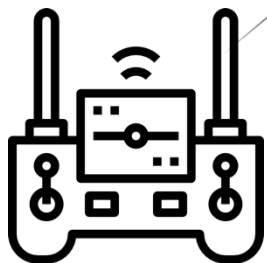
Count = 15



UAV

Mission_item_int(5)

MAV LINK



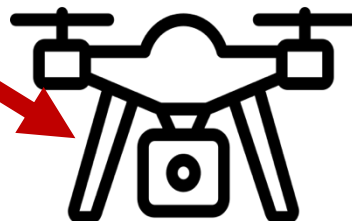
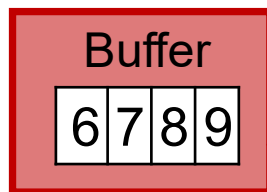
GCS

Widely Used Protocols Are Susceptible to Stealthy Attacks

Resource Usage
Exploit¹

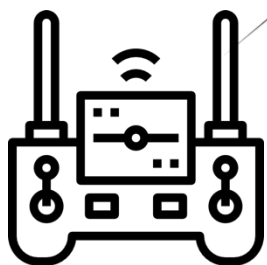
Mission 2:

Count = 15



UAV

UAV starts behaving as mission 1
in the middle of mission 2



GCS

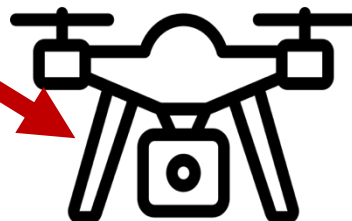
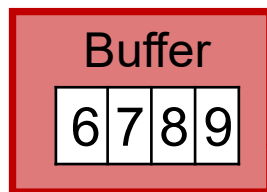


Existing Defenses Miss the Logical Layer

Resource Usage
Exploit¹

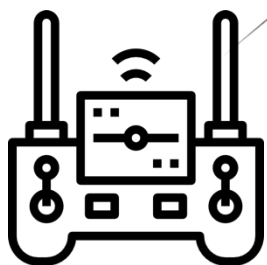
Mission 2:

Count = 15



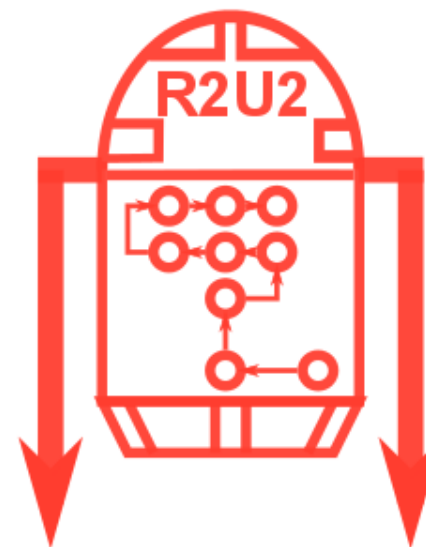
UAV

UAV starts behaving as mission 1
in the middle of mission 2



GCS

Problem: A stealthy attack uses
the protocol exactly as specified
to induce undesired behavior,
which makes it invisible to the
defenses we usually rely on



Intrusion detection systems:
No anomalous behavior or
virus signature

Physics based anomaly
detection: no anomalous
physical deviation.

DATUM: Proof of Concept, Two Blockers to Deployment

Global Refined Multiparty Session Types (GRMPST)

```
GCS → UAV: MISSION COUNT( $n\{1 \leq n < \text{LIMIT}\}$ )  
 $\mu.T(\text{curr}\{0 \leq \text{curr} < n\})(\text{curr} = 0)$   
UAV → GCS:  
  MISSION REQUEST INT( $\text{req}\{\text{req} = \text{curr} < n\}$ )  
  MISSION ITEM INT( $\text{item}\{\text{item} = \text{curr} < n\}$ )  
  T( $\text{curr} + 1$ )  
⊕ MISSION ACK( $\text{ack}\{\text{type} = \text{ERROR} \vee \text{curr} = n\}$ ).End
```

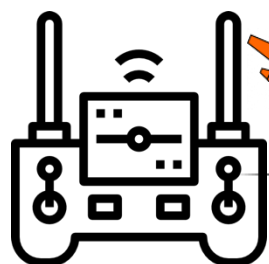
F* Specification



compile

Compile time

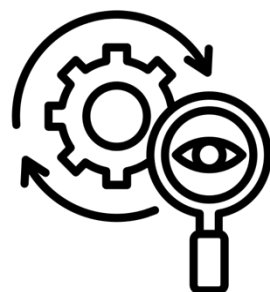
Runtime



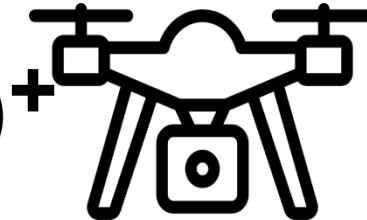
GCS



MAVLINK



Datum



UAV

DATUM: Proof of Concept, Two Blockers to Deployment

Global Refined Multiparty Session Types (GRMPST)

```
GCS → UAV: MISSION COUNT( $n\{1 \leq n < \text{LIMIT}\}$ )  
 $\mu.T(\text{curr}\{0 \leq \text{curr} < n\})(\text{curr} = 0)$   
UAV → GCS:  
  MISSION REQUEST INT( $\text{req}\{\text{req} = \text{curr} < n\}$ )  
  MISSION ITEM INT( $\text{item}\{\text{item} = \text{curr} < n\}$ )  
  T( $\text{curr} + 1$ )  
⊕ MISSION ACK( $\text{ack}\{\text{type} = \text{ERROR} \vee \text{curr} = n\}$ ).End
```

F* Specification



compile

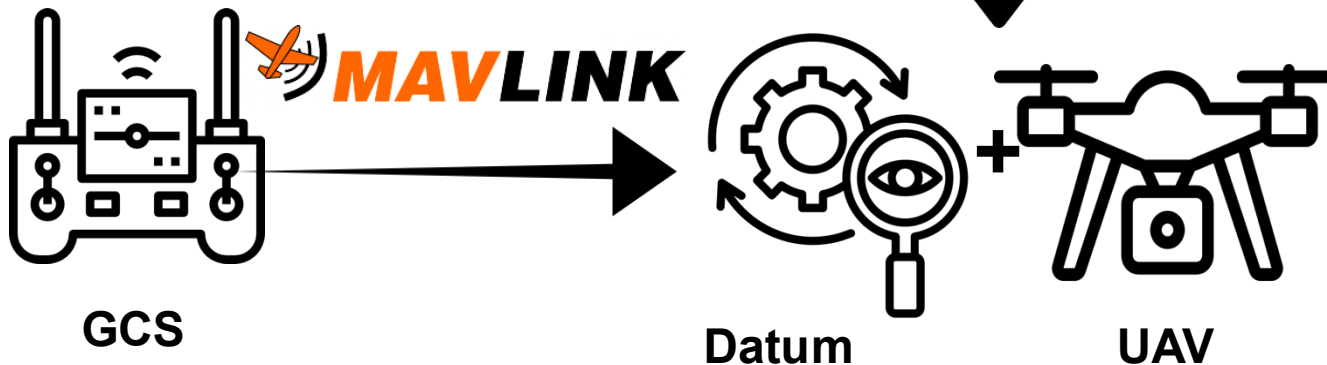
Limitation #1: Specification Burden

Reason: Extrinsic Verification

```
(#nparticipants : nat) ->  
(#rec_contexts : list (session_type u))  
(from : nat) -> (to : nat) ->  
(1 : (list (label * (dtuple4 ...
```

Extra proof arguments

Compile time
Runtime



DATUM: Proof of Concept, Two Blockers to Deployment

Global Refined Multiparty Session Types (GRMPST)

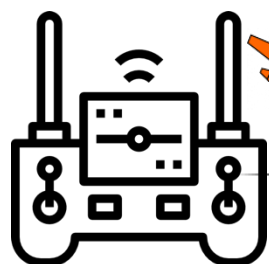
```
GCS → UAV: MISSION COUNT( $n\{1 \leq n < \text{LIMIT}\}$ )  
 $\mu.T(\text{curr}\{0 \leq \text{curr} < n\})(\text{curr} = 0)$   
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  MISSION ITEM INT( $\text{item}\{\text{item} = \text{curr} < n\}$ )  
  T( $\text{curr} + 1$ )  
⊕ MISSION ACK( $\text{ack}\{\text{type} = \text{ERROR} \vee \text{curr} = n\}$ ).End
```

F* Specification

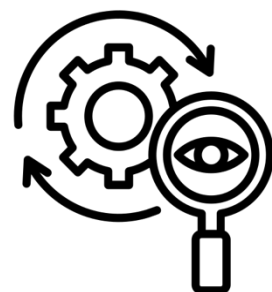


compile

Compile time
Runtime



GCS



Datum

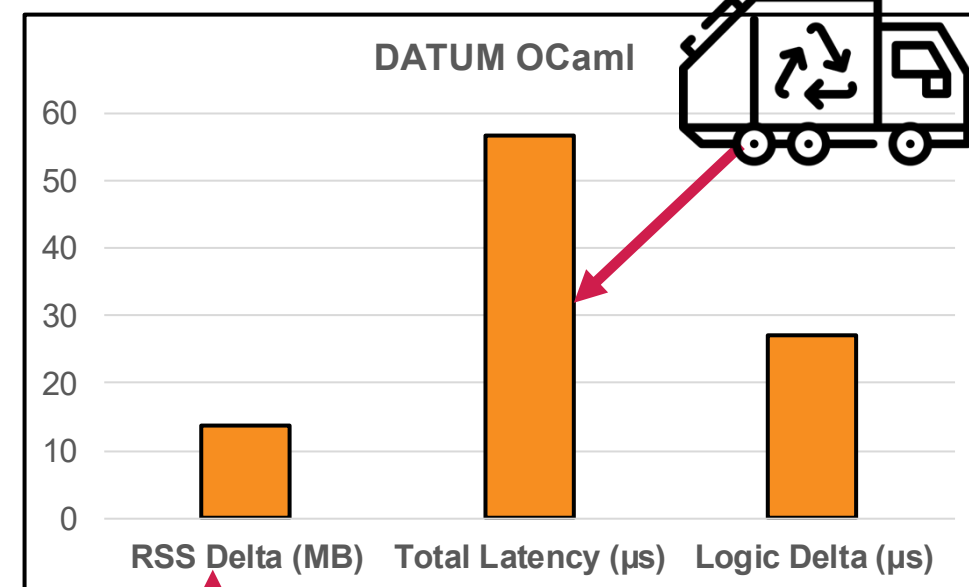


UAV

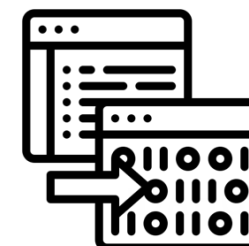
Limitation #2 : The Deployment Gap

Reason: OCaml runtime

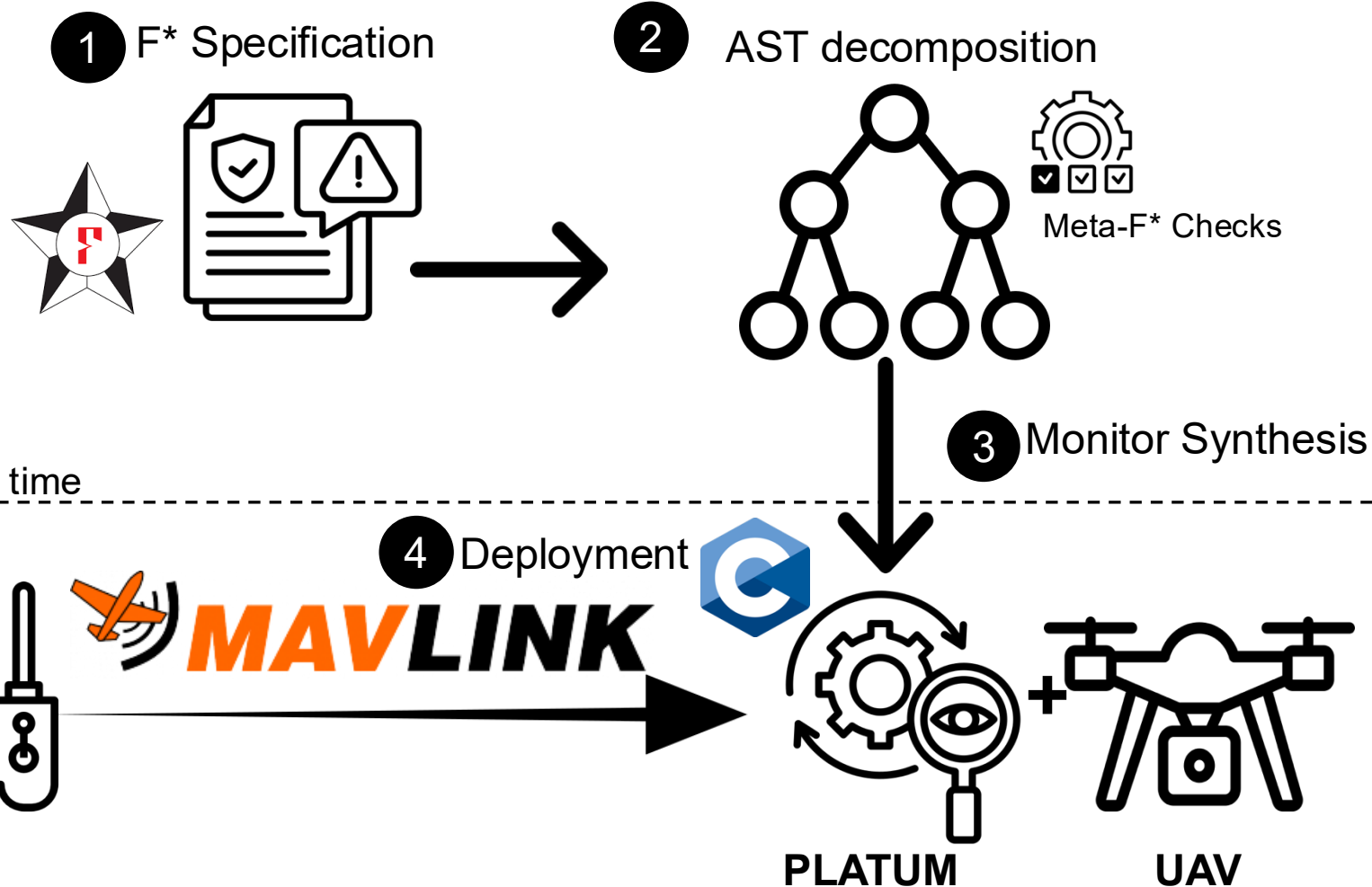
Garbage Collector



OCaml
Compiler



Platum: *Five Inputs In. Certified C Monitor Out*



Contributions

1- A minimal DSL: session inputs, no proof scaffolding, and automated checks.

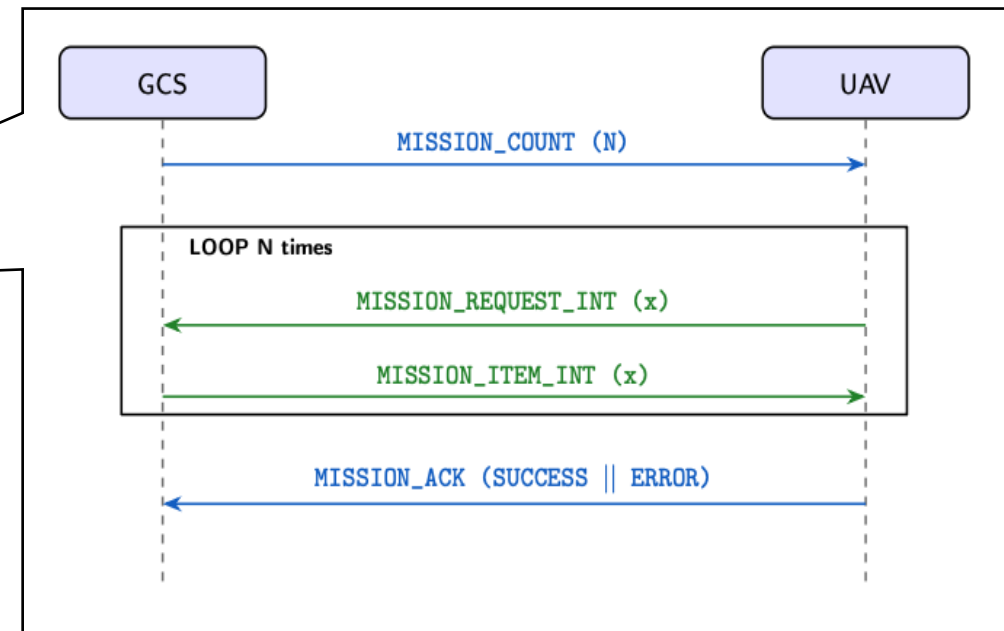
2- Synthesis of performant monitors from wellformed session types.

3- 4x latency and 40% memory reduction vs. DATUM under ArduPilot SITL.

GRMPSTs Enforce Ordering, Preconditions and Value Constraints at the Network Boundary

Global Refined Multiparty Session Types (GRMPST)

```
GCS → UAV: MISSION COUNT( $n\{1 \leq n < \text{LIMIT}\}$ )  
 $\mu.T(\text{curr}\{0 \leq \text{curr} < n\})(\text{curr} = 0)$   
UAV → GCS:  
  MISSION REQUEST INT( $\text{req}\{\text{req} = \text{curr} < n\}$ )  
  MISSION ITEM INT( $\text{item}\{\text{item} = \text{curr} < n\}$ )  
   $T(\text{curr} + 1)$   
 $\oplus \text{MISSION ACK}(\text{ack}\{\text{type} = \text{ERROR} \vee \text{curr} = n\}).\text{End}$ 
```



GRMPSTs Enforce Ordering, Preconditions and Value Constraints at the Network Boundary

Global Refined Multiparty Session Types (GRMPST)

GCS $\rightarrow$ **UAV**: MISSION COUNT($n\{1 \leq n < \text{LIMIT}\}$)

1 $\mu.T(\text{curr}\{0 \leq \text{curr} < n\})(\text{curr} = 0)$

UAV $\rightarrow$ **GCS**:

MISSION REQUEST INT($\text{req}\{\text{req} = \text{curr} < n\}$)

MISSION ITEM INT($\text{item}\{\text{item} = \text{curr} < n\}$)

$T(\text{curr} + 1)$

3 $\oplus$ MISSION ACK($\text{ack}\{\text{type} = \text{ERROR} \vee \text{curr} = n\}$).End

- 1 Create a recursive variable *curr* that must be less than or equal to *n*
- 2 Ensure messages are uploaded in the correct order and not skipped
- 3 The protocol must only complete with *curr* = *n* or with an error message.

GRMPSTs Enforce Ordering, Preconditions and Value Constraints at the Network Boundary

Resource Usage
Exploit¹

Mission 1:

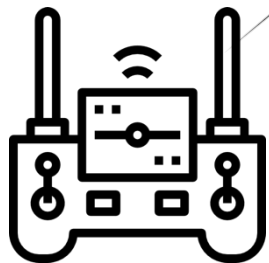
Count = 10

Buffer

6	7	8	9
---	---	---	---



UAV



GCS

2

Ensure messages are uploaded in the correct order and not skipped

GRMPSTs Enforce Ordering, Preconditions and Value Constraints at the Network Boundary

Resource Usage
Exploit¹

Mission 1:

Count = 10

Buffer

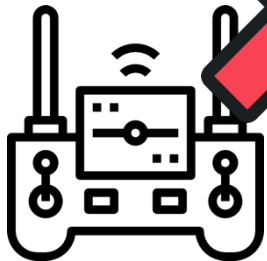
6	7	8	9
---	---	---	---



UAV

Mission_ack(Success)

MAV LINK



GCS

3

The protocol must only complete with *curr* = *n* or with an error message.

Extrinsic vs. Intrinsic — The DSL Design Choice

1 F* Specification



5 key arguments for a GRMPST

Diagram illustrating the 5 key arguments for a GRMPST (Goal Refinement Method for Policy Specification and Testing):

1- sender: *GCS*

2- receiver: *UAV*

3- label: *MISSION COUNT*

4- variable: *n*

5- refinement: *mission count {1 ≤ count(n) < LIMIT}*

The diagram shows the expression: *GCS* → *UAV*: *MISSION COUNT*(*n*: *mission count* {1 ≤ count(*n*) < *LIMIT*})

Extrinsic vs. Intrinsic — The DSL Design Choice

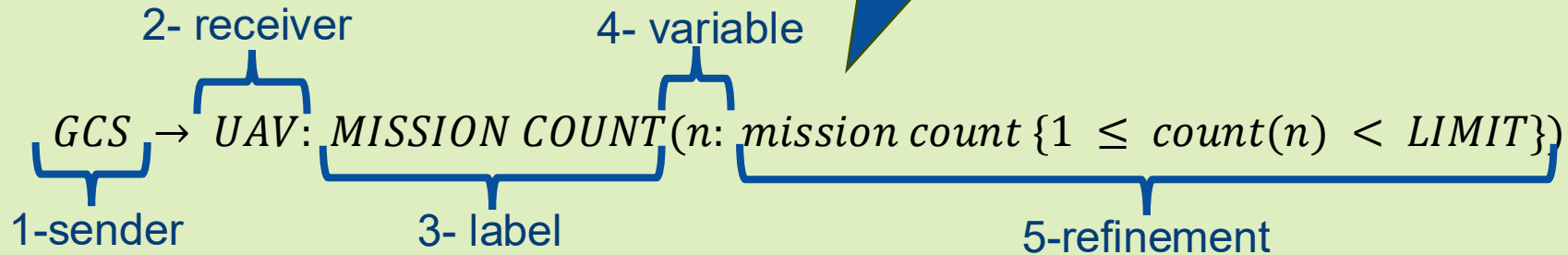
Must check these properties

- 1- Label Uniqueness
- 2- Guarded Recursion
- 3- Session Fidelity
- 4- Global Progress

1 F* Specification



5 key arguments for a GRMPST



Extrinsic vs. Intrinsic — The DSL Design Choice

Intrinsic Verification

Type checks = is well-formed

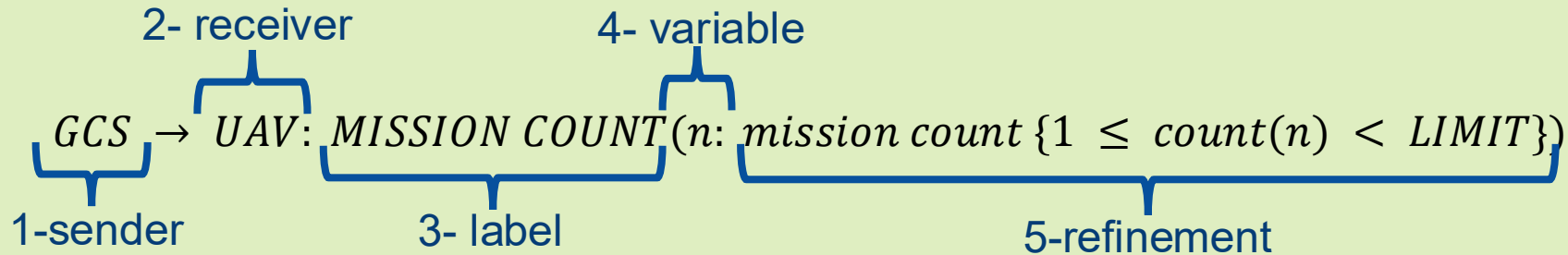
Pros: Cannot construct/compile a malformed term

Cons: Must provide the type checker in enough information to check well-formedness (extra inputs)

1 F* Specification



5 key arguments for a GRMPST



Extrinsic Verification

Type checks $\neq$ is well-formed (checked separately)

Pros: Minimal syntactic surface (only required inputs)

Cons: Must ensure checks happen at some point.

Extrinsic vs. Intrinsic — The DSL Design Choice

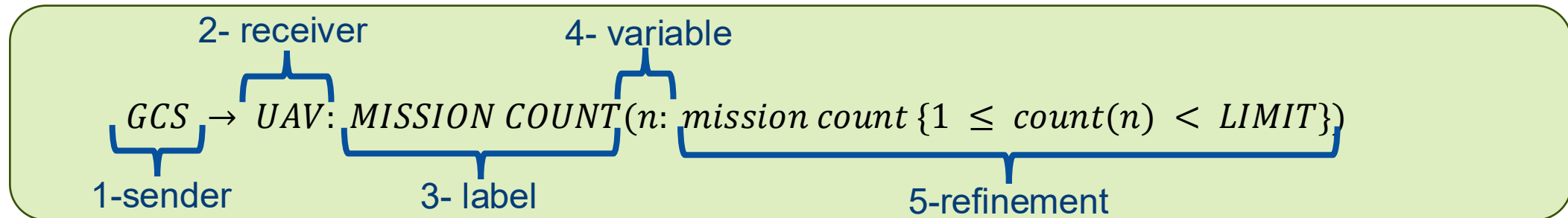
OLD - Intrinsic verification (Datum)

```
(#nparticipants : nat) ->  
(#rec_contexts : list (session_type u)) ->  
(from : nat) -> (to : nat) ->  
(l : (list (label * (dtuple4 ...
```

1 F* Specification



5 key arguments for a GRMPST



Extrinsic Verification

Type checks $\neq$ is well-formed (checked separately)

Pros: Minimal syntactic surface (only required inputs)

Cons: Must ensure checks happen at some point.

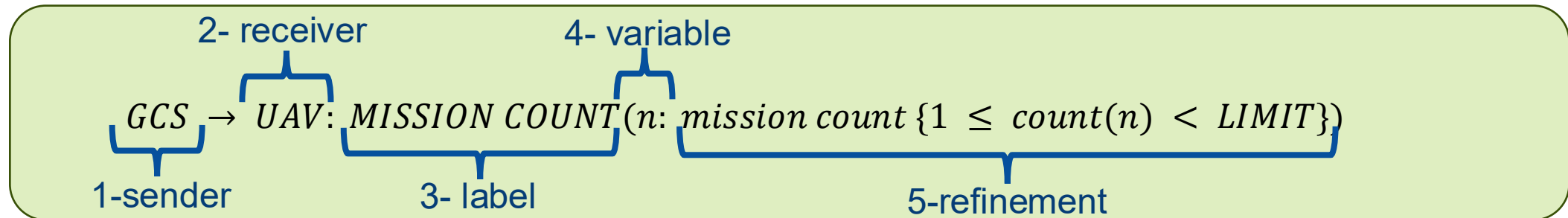
Extrinsic vs. Intrinsic — The DSL Design Choice

OLD - Intrinsic verification (Datum)

```
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(#rec_contexts : list (session_type u)) ->  
(from : nat) -> (to : nat) ->  
(l : (list (label * (dtuple4 ...
```

Extra proof inputs

5 key arguments for a GRMPST



Extrinsic Verification

Type checks $\neq$ is well-formed (checked separately)

Pros: Minimal syntactic surface (only required inputs)

Cons: Must ensure checks happen at some point.

1 F* Specification



Extrinsic vs. Intrinsic — The DSL Design Choice

OLD - Intrinsic verification (Datum)

```
(#nparticipants : nat) ->  
(#rec_contexts : list (session_type u)) ->  
(from : nat) -> (to : nat) ->  
(l : (list (label * (dtuple4 ...
```

Extra proof inputs

5 key arguments for a GRMPST

2- receiver
4- variable
1- sender
3- label
5- refinement

$GCS \rightarrow UAV: MISSION_COUNT(n: mission_count \{1 \leq count(n) < LIMIT\})$

NEW - Intrinsic Verification (Platum)

2- receiver
4- variable
1- sender
3- label
5- refinement

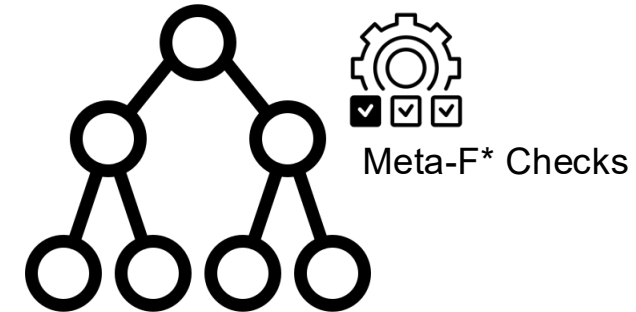
$(0 \dashrightarrow 1) [Option\ "MISSION_COUNT" (\lambda\ (cnt : mission_count \{ cnt \geq 1 \ \&\& \ cnt < mission_limit \})$

1 F* Specification



Meta-F* Decomposes the AST Into an Inspectable Protocol Graph

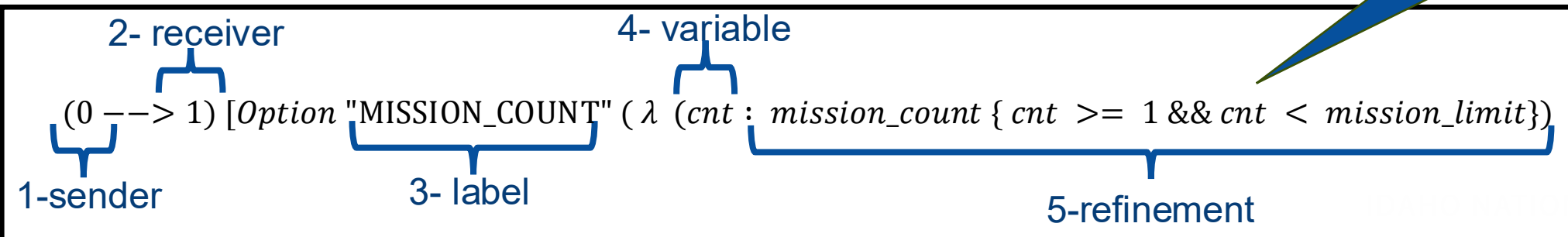
2 AST decomposition



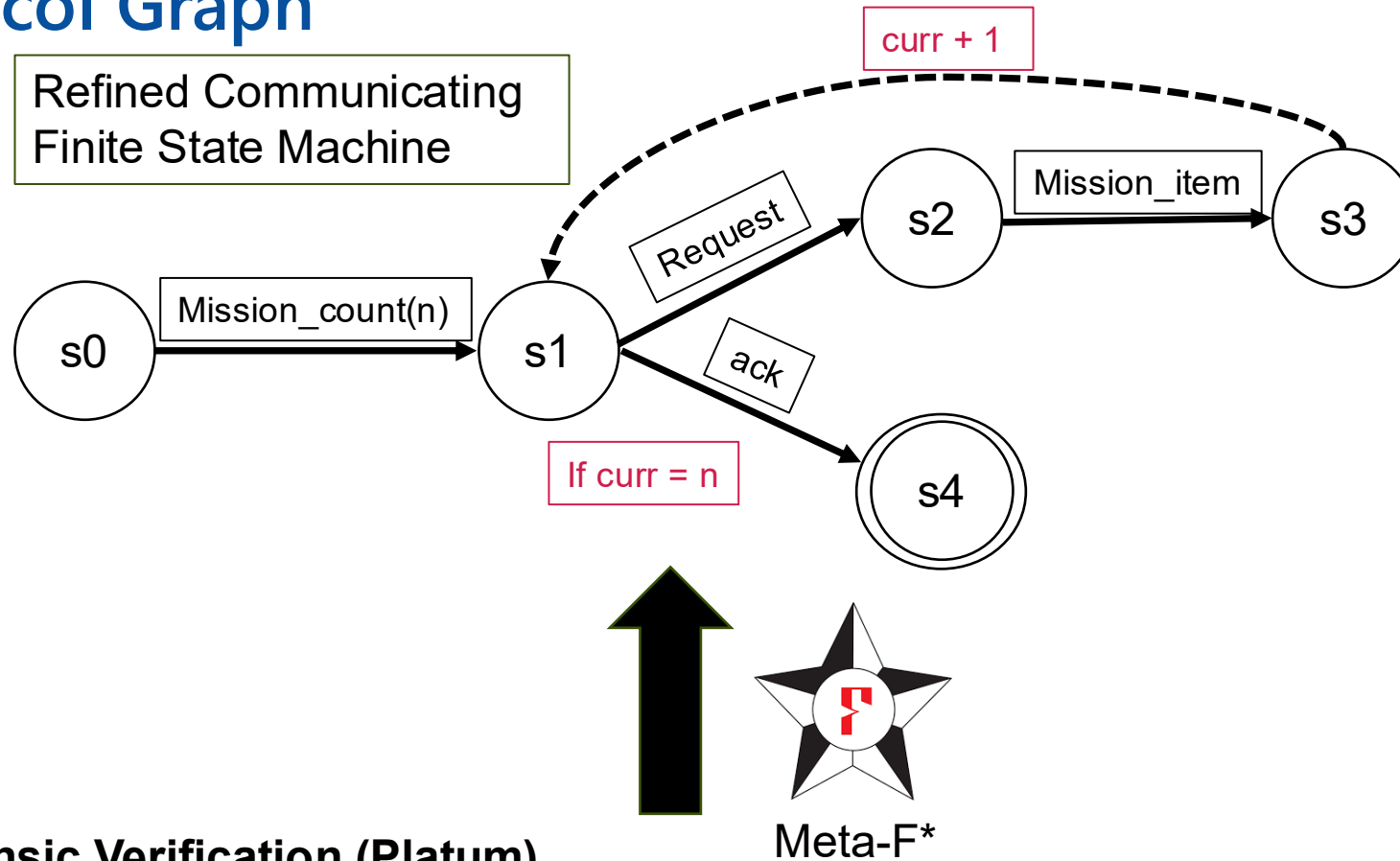
Must check these properties

- 1- Label Uniqueness
- 2- Guarded Recursion
- 3- Session Fidelity
- 4- Global Progress

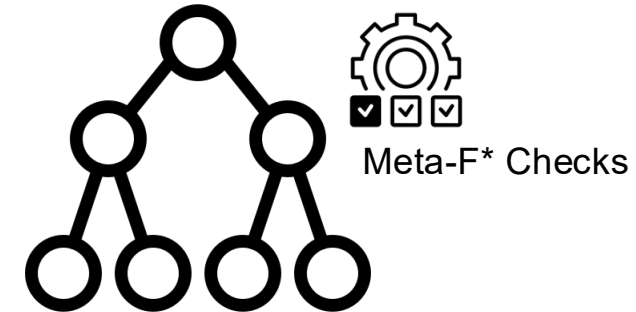
NEW - Intrinsic Verification (Platum)



Meta-F* Decomposes the AST Into an Inspectable Protocol Graph



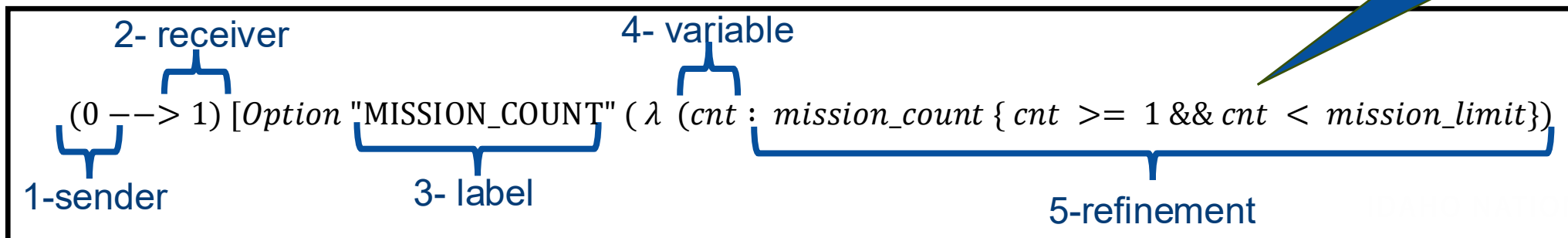
2 AST decomposition



Must check these properties

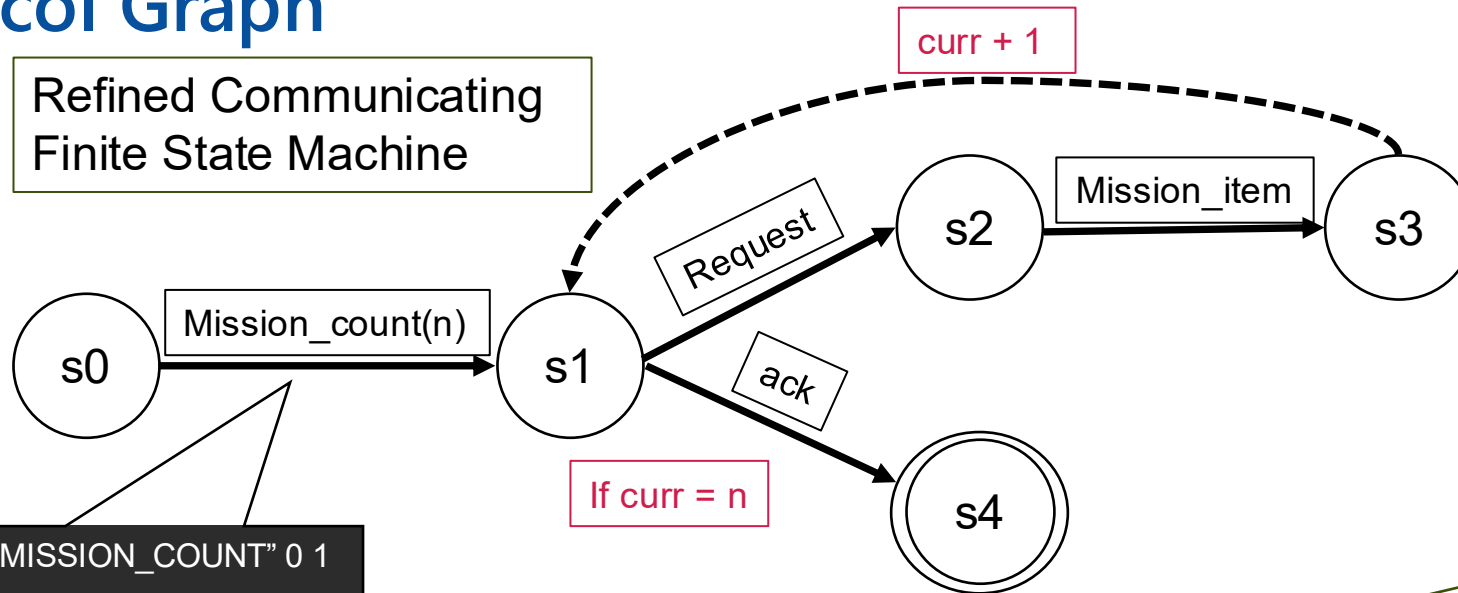
- 1- Label Uniqueness
- 2- Guarded Recursion
- 3- Session Fidelity
- 4- Global Progress

NEW - Intrinsic Verification (Platum)



Meta-F* Decomposes the AST Into an Inspectable Protocol Graph

Refined Communicating
Finite State Machine

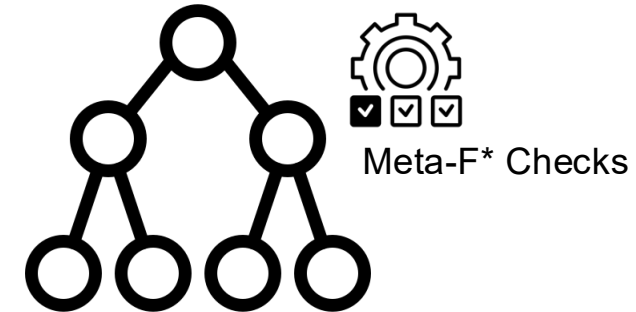


```
Mkprotocol_transition "MISSION_COUNT" 0 1
[
  Mkvar "mc_target_system" KindInt;
  Mkvar "mc_target_component" KindInt;
  Mkvar "mc_count" KindInt;
  Mkvar "mc_mission_type" KindInt;
  Mkvar "mc_opaque_id" KindInt
]
(GAnd (GGe (TVar "mc_count") (TIntLit
1))) (GLt (TVar "mc_count") (TIntLit 65535)))
```

```
let rec has_duplicates (l: list label) (seen: list label)
  Tot bool =
  match l with
  | [] → false
  | hd :: tl →
    if L.mem hd seen && hd <> "RECUR" then true
    else has_duplicates tl (hd :: seen)

let check_unique_labels (labels: list label) : bool =
  not (has_duplicates labels [])
```

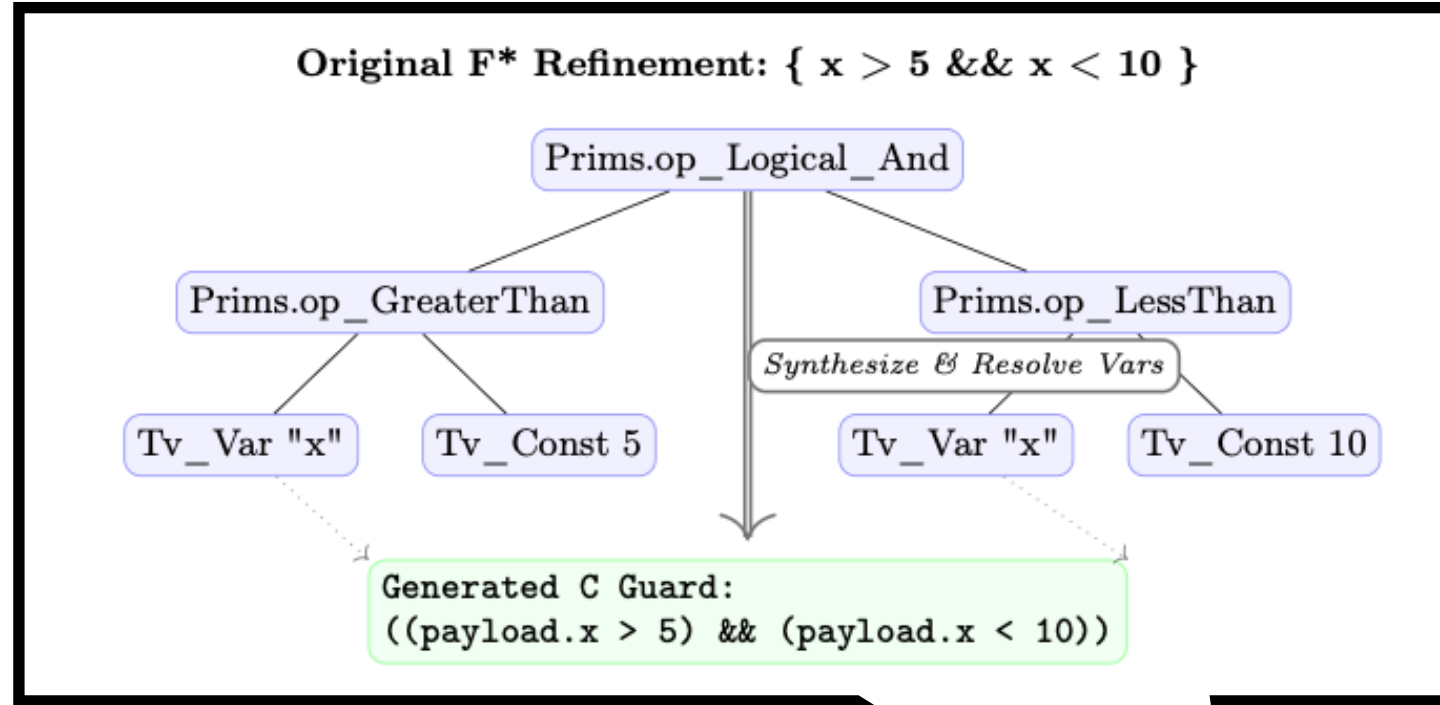
2 AST decomposition



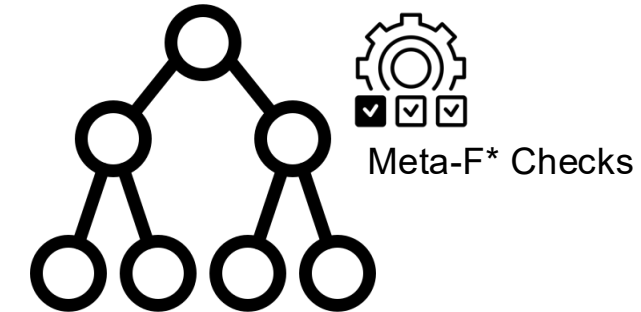
Must check these properties

- 1- Label Uniqueness
- 2- Guarded Recursion
- 3- Session Fidelity
- 4- Global Progress

Type Erasure Without Losing the Constraints



2 AST decomposition



NEW - Intrinsic Verification (Platum)

$(0 \dashrightarrow 1) [Option \text{ "MISSION_COUNT" } (\lambda \text{ (cnt : mission_count } \{ cnt \geq 1 \ \&\& \ cnt < mission_limit \})$

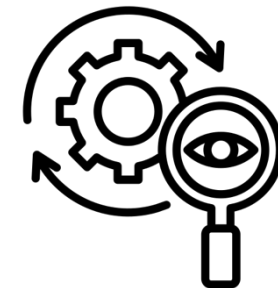
5-refinement

Type Erasure Without Losing the Constraints

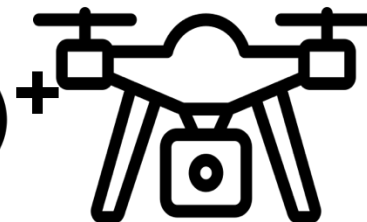
Advantages:

- 1- Allocation-free
- 2- $O(1)$ state lookup (jump table)
- 3- Bounded WCET (no GC pauses)

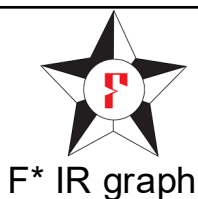
3 Monitor Synthesis



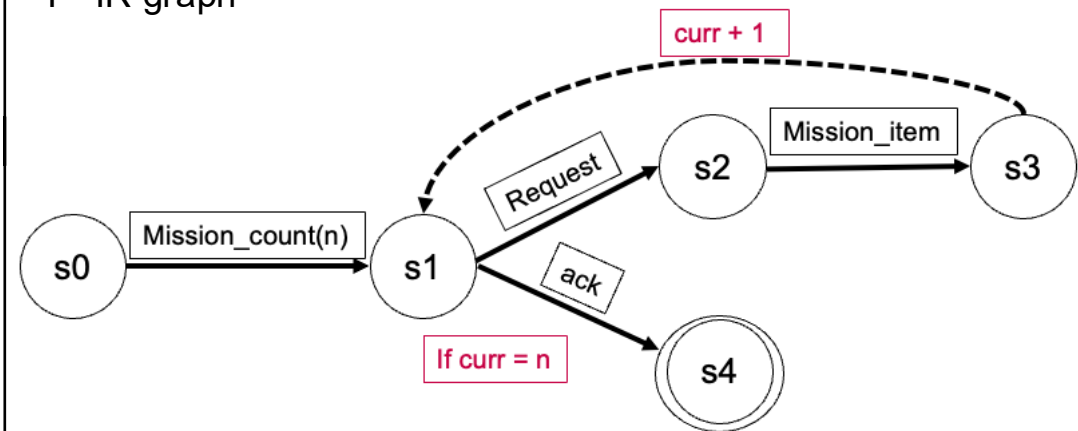
PLATUM



UAV



F* IR graph



Synthesis

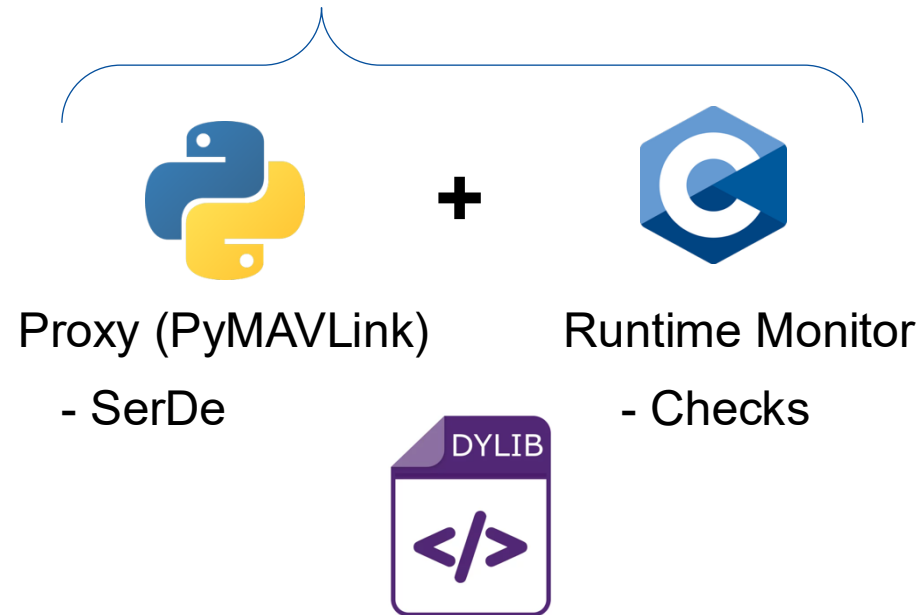
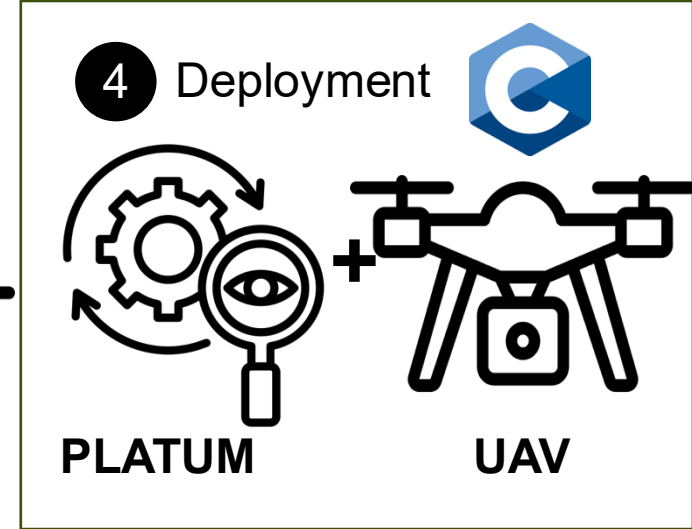
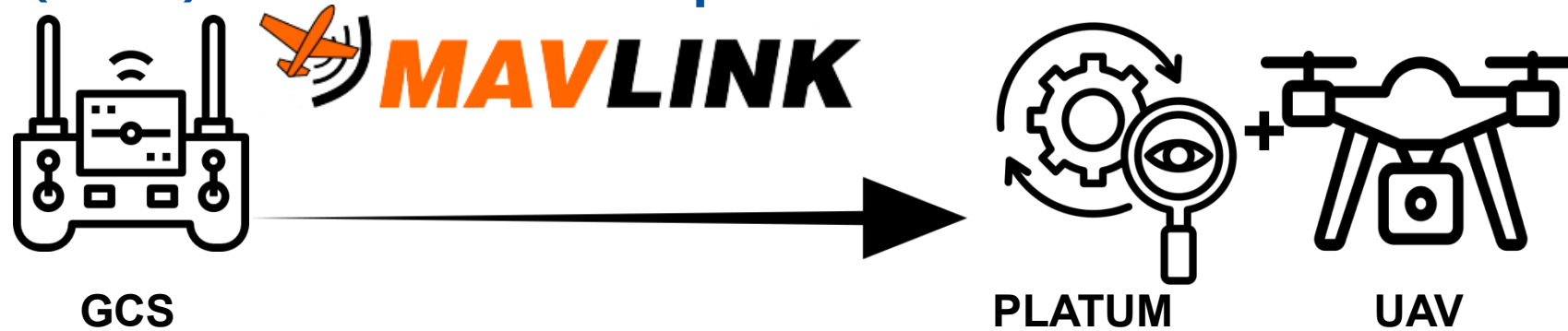
```
bool monitor_step(monitor_t* mon, const mavlink_message_t*
msg) {
    switch (mon->state) {
        // State 0: Expecting MISSION_COUNT
        // GCS -> UAV: MISSION_COUNT(N: Z {N >= 1 /\ N < LIMIT})
        case STATE_IDLE_0: {
            if (msg->msgid == MAVLINK_MSG_ID_MISSION_COUNT) {
                mavlink_mission_count_t payload;
                mavlink_msg_mission_count_decode(msg, &payload);
                // Check Refinement: {cnt >= 1 && cnt < mission_item_limit}
                if (payload.count >= 1 && payload.count < 65535) {
                    // Update Context (F* bound variables)
                    mon->ctx.cnt = payload.count;
                    mon->ctx.curr = 0; // Initialize Mu variable
                }
                // Transition
                mon->state = STATE_LOOP_REQ_1;
                return true;
            }
        }
    }
}
```



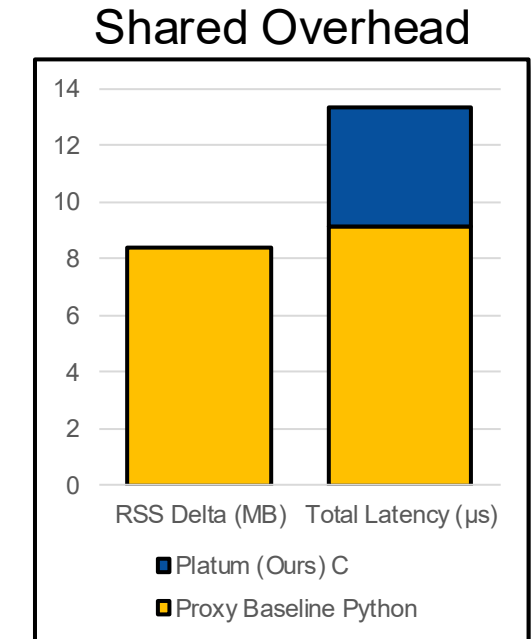
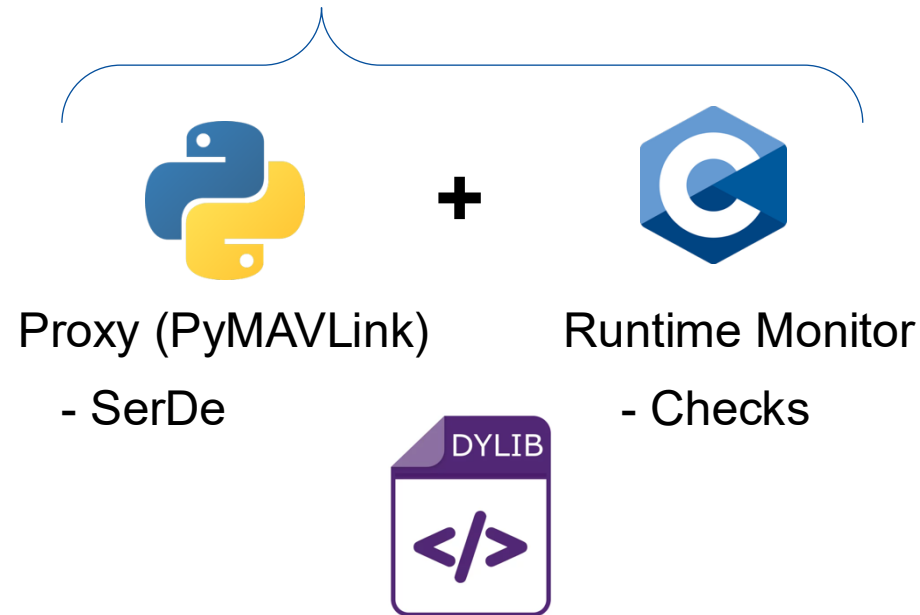
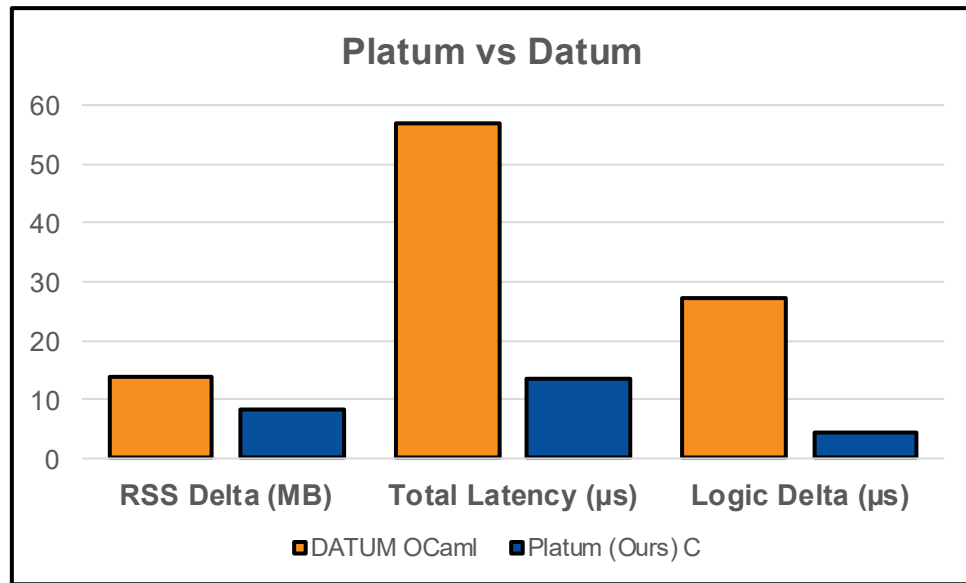
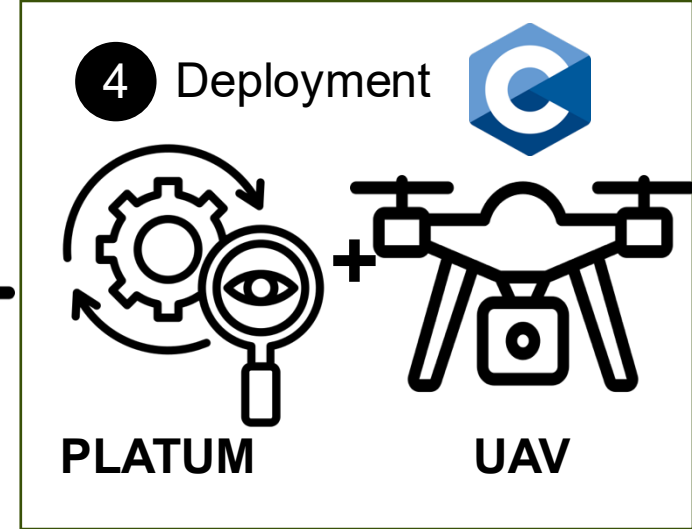
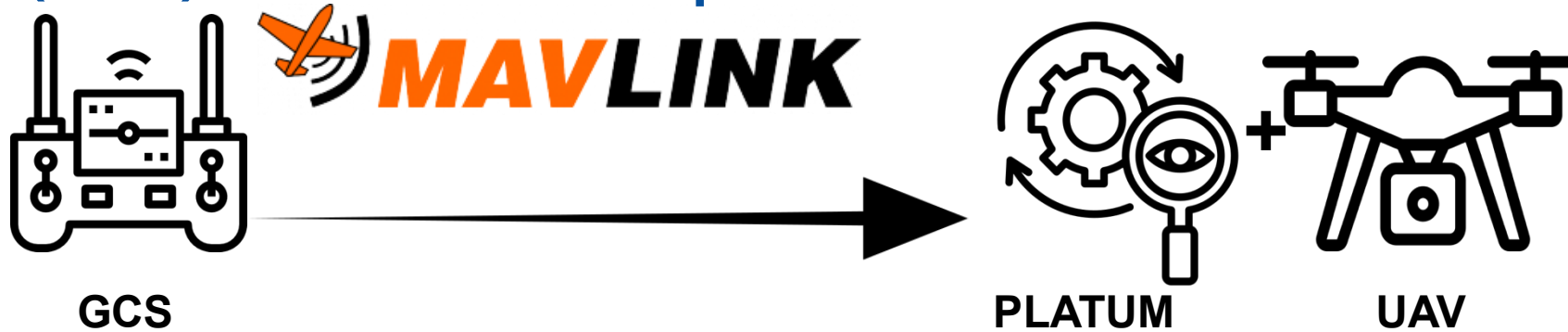
CFSM

Must pass all 4 wellformedness checks

Platum decreases Latency (4x) and Memory (40%) Overhead Compared to Datum



Platum decreases Latency (4x) and Memory (40%) Overhead Compared to Datum

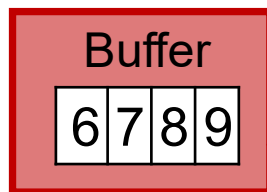


Platum: *Five Inputs In. Certified C Monitor Out*

Resource Usage
Exploit¹

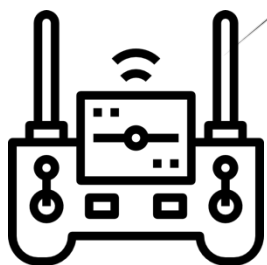
Mission 2:

Count = 15



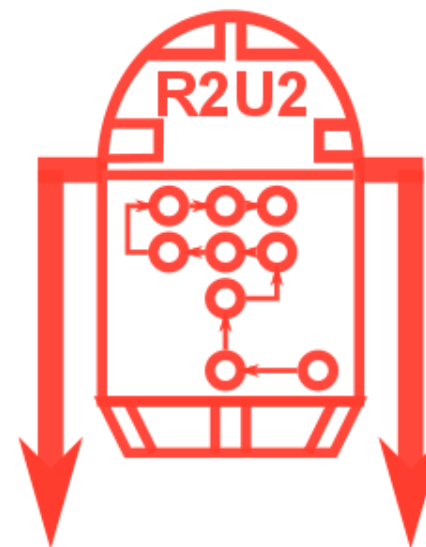
UAV

UAV starts behaving as mission 1
in the middle of mission 2



GCS

Problem: A stealthy attack uses
the protocol exactly as specified
to induce undesired behavior,
which makes it invisible to the
defenses we usually rely on

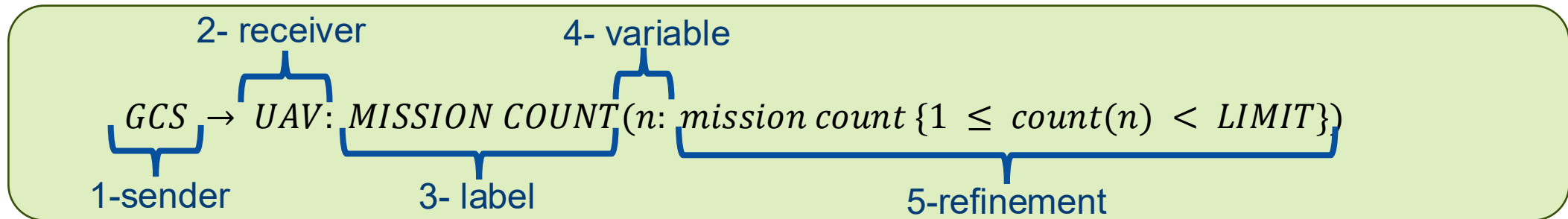


Intrusion detection systems:
No anomalous behavior or
virus signature

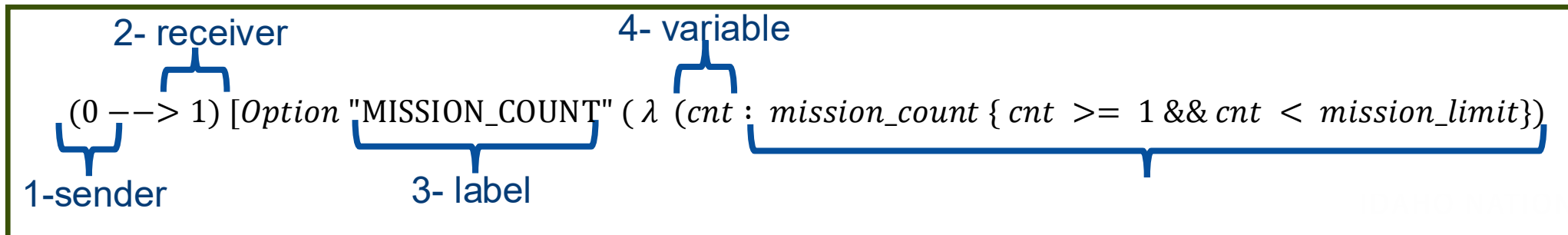
Physics based anomaly
detection: no anomalous
physical deviation.

Platum: *Five Inputs In. Certified C Monitor Out*

1 F* Specification



NEW - Intrinsic Verification (Platum)

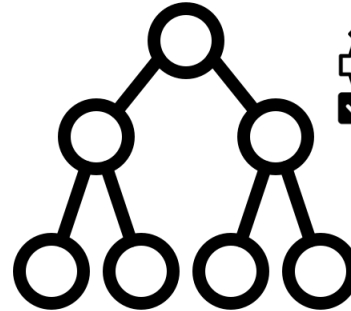


Platum: *Five Inputs In. Certified C Monitor Out*

1 F* Specification



2 AST decomposition

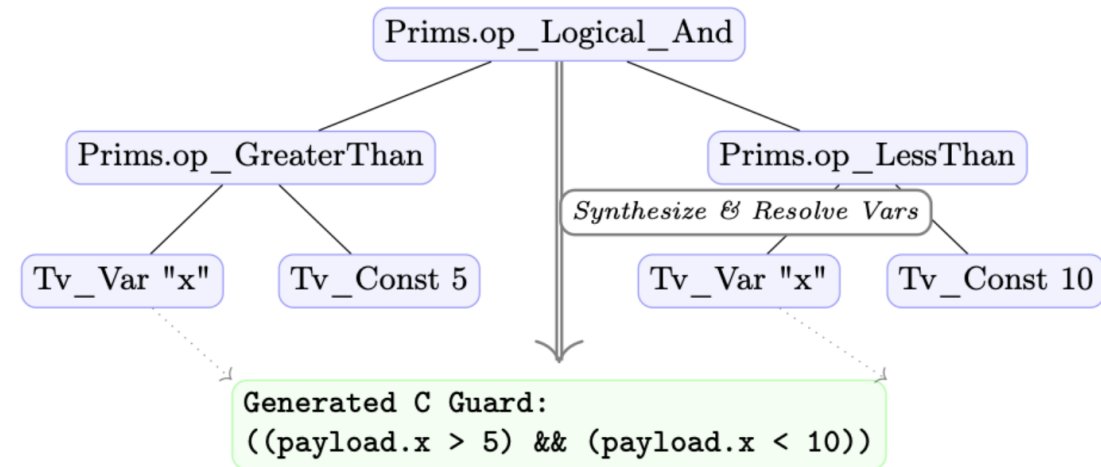


Meta-F* Checks

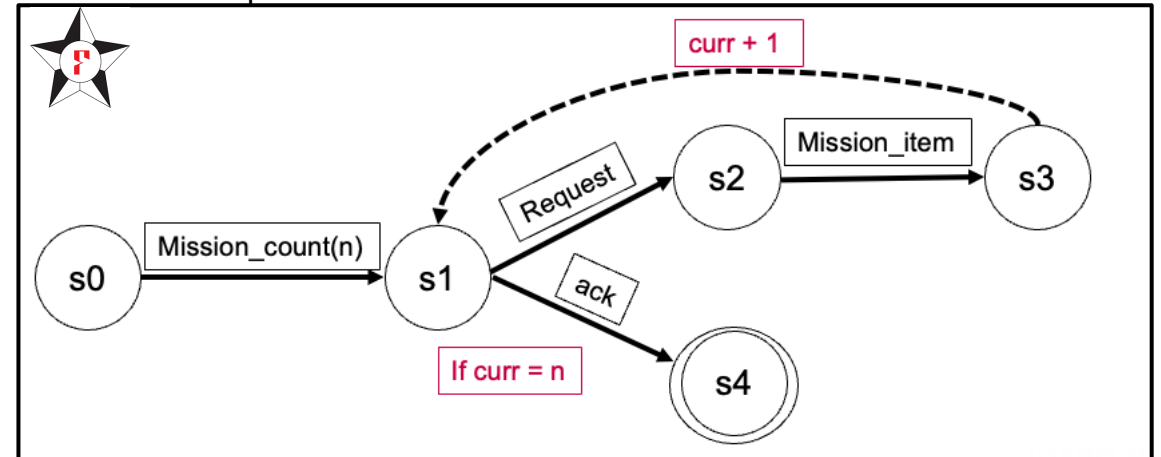
Wellformedness Checks

- 1- Label Uniqueness
- 2- Guarded Recursion
- 3- Session Fidelity
- 4- Global Progress

Original F* Refinement: $\{ x > 5 \ \&\& \ x < 10 \}$



Intermediate representation



Platum: *Five Inputs In. Certified C Monitor Out*

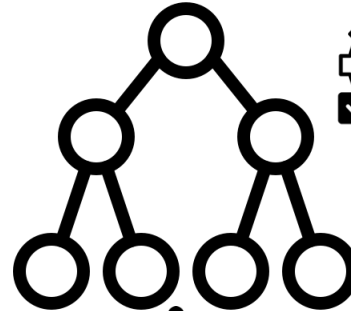


C FSM

1 F* Specification

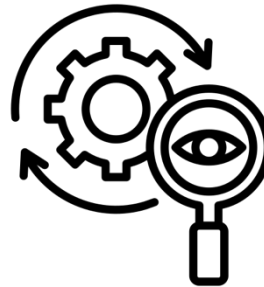


2 AST decomposition

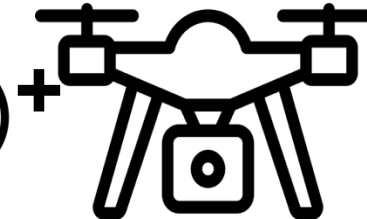


Meta-F* Checks

3 Monitor Synthesis



PLATUM



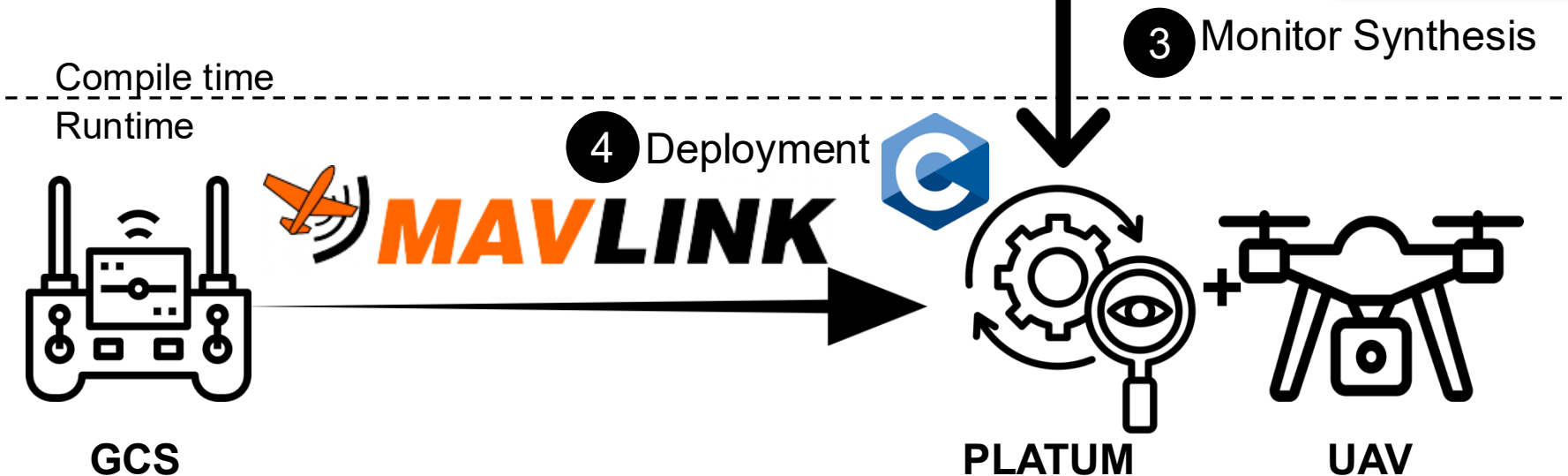
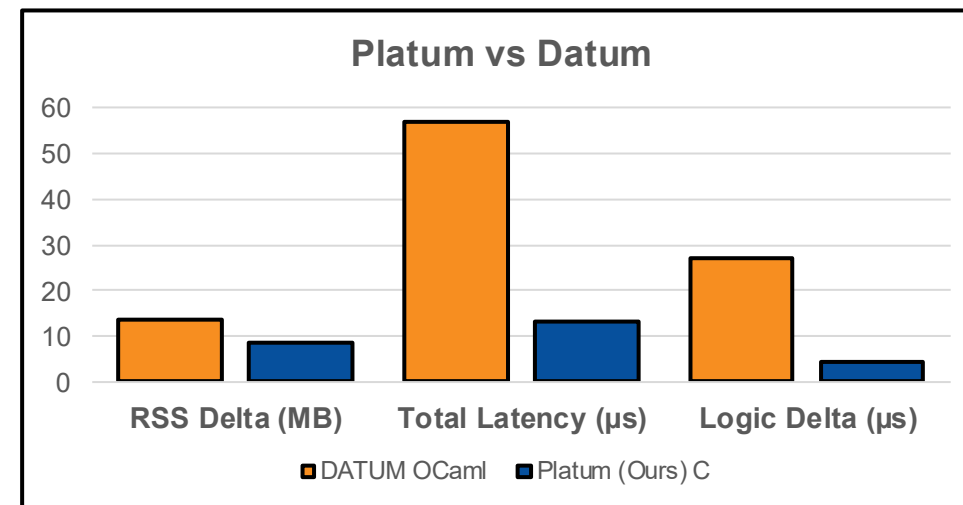
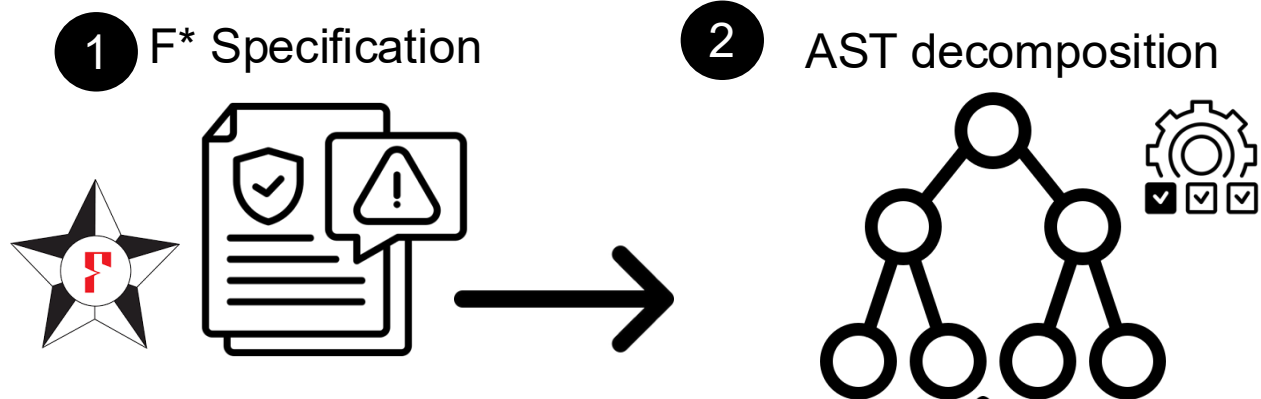
UAV

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msg) {
    switch (mon->state) {
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                    // Update Context (F* bound variables)
                    mon->ctx.cnt = payload.count;
                    mon->ctx.curr = 0; // Initialize Mu variable
                }
                // Transition
                mon->state = STATE_LOOP_REQ_1;
                return true;
            }
        }
    }
}
```

Advantages:

- 1- Allocation-free
- 2- $O(1)$ state lookup (jump table)
- 3- Bounded WCET (no GC pauses)

Platum: *Five Inputs In. Certified C Monitor Out*

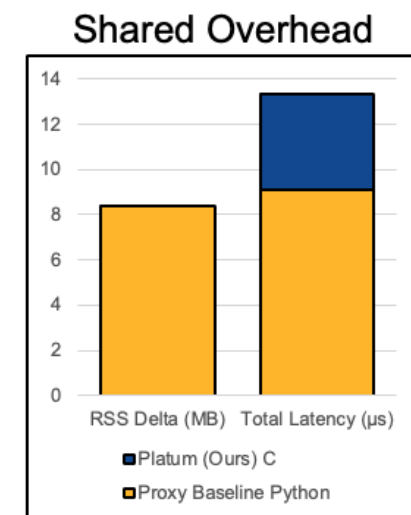
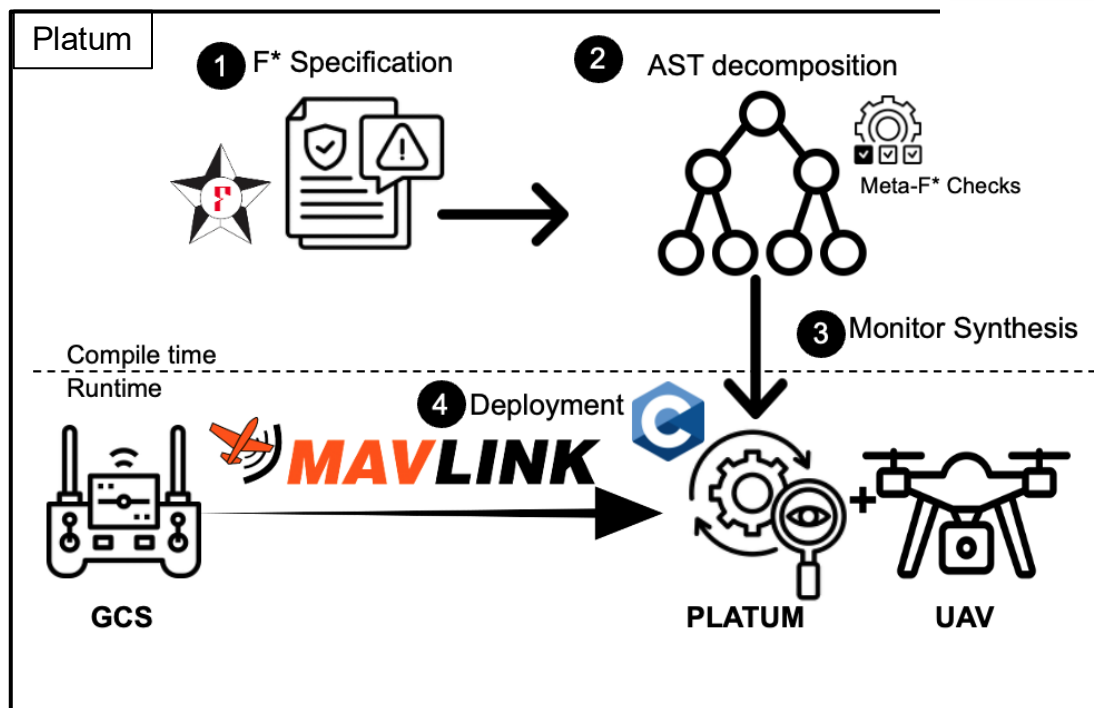
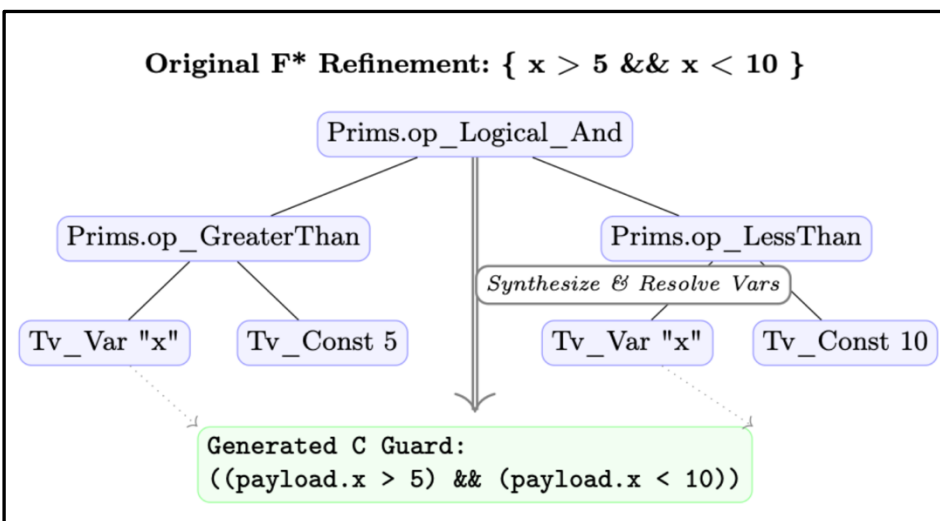
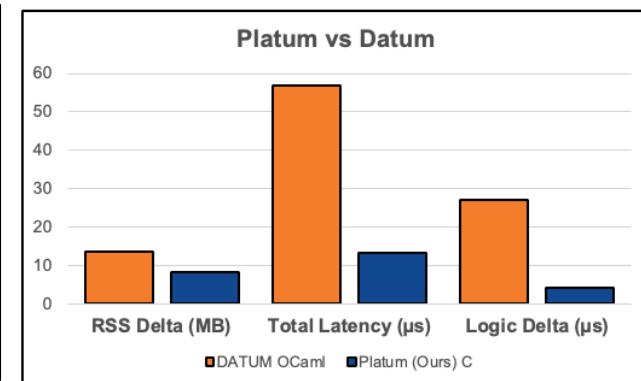
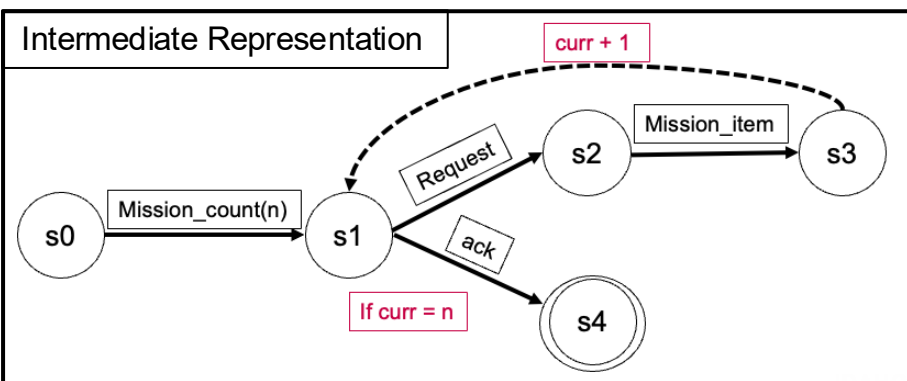


Five inputs in, verified C monitor out — 4x faster and 40% leaner than the state of the art.

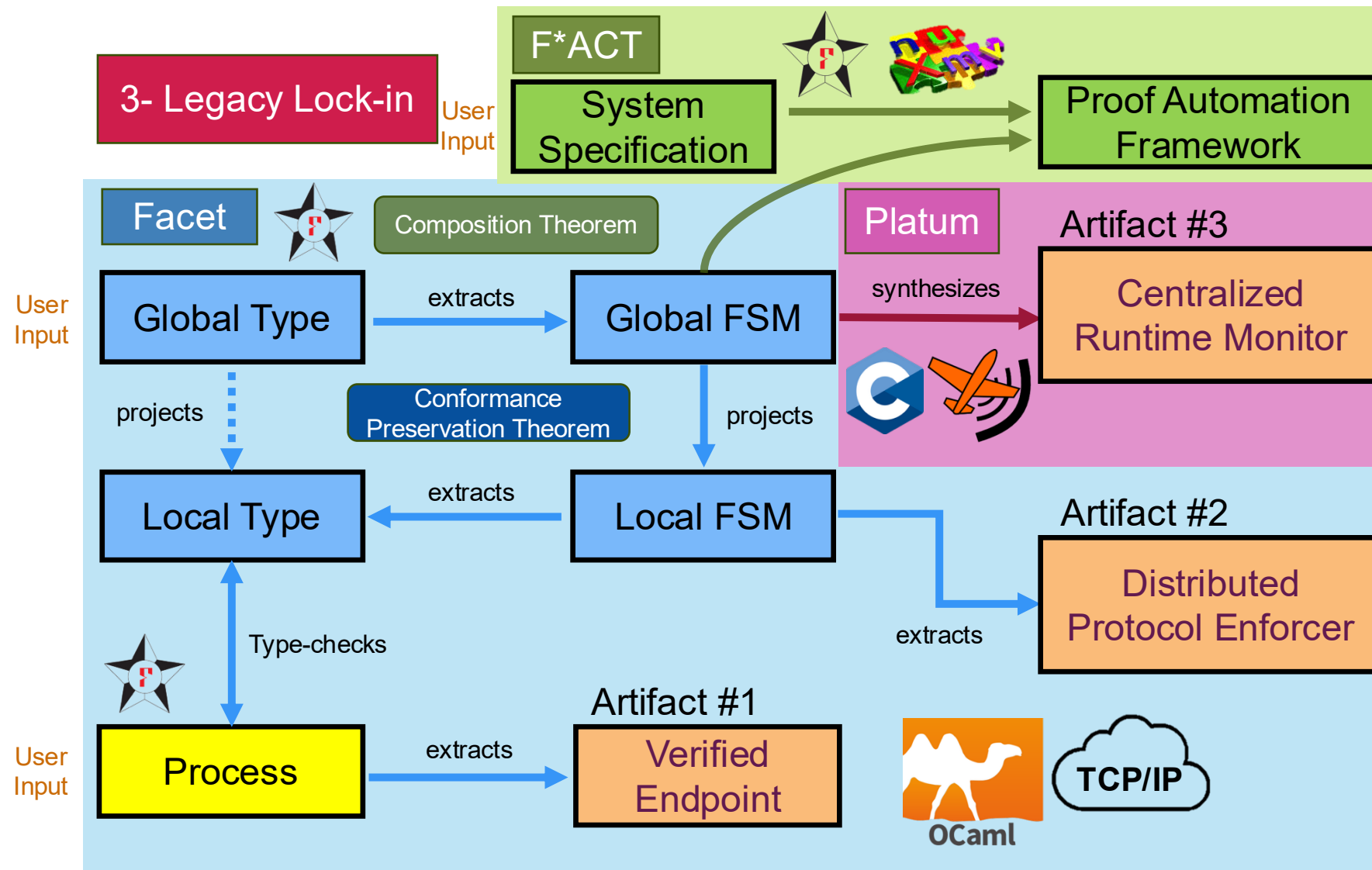
OPEN TO WORK



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arthur.amorim@inl.gov
Website



Platum: Retrofitting Legacy Protocols.



Evaluation numbers

System	Language	RSS Delta (MB)	Total Latency (μ s)	Logic Delta (μ s)
DATUM	OCaml	13.72	56.74	27.14
Proxy Baseline	Python	8.39	9.11	0
Platum (Ours)	C	8.39	13.33	4.22